

Hide in Plain Sight: Enabling Mobile Applications and Data Analytics with Local Differential Privacy

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Location data collected from individual devices
(Source: New York Times 12/2018)

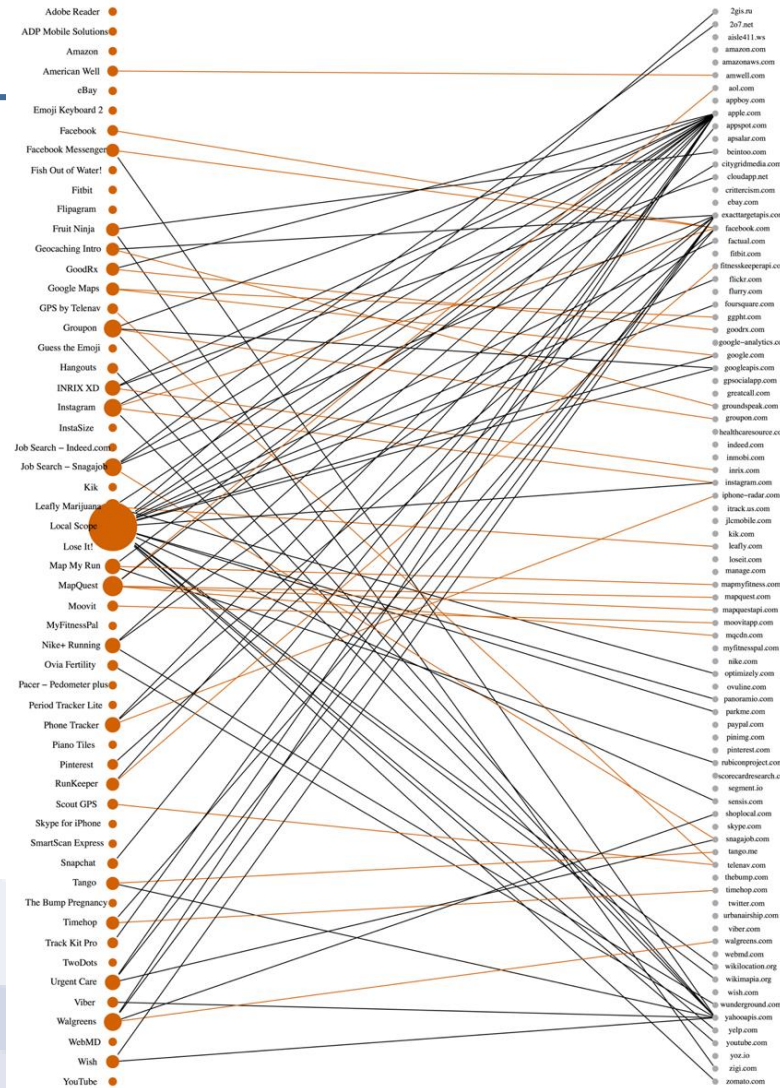




Over 235 million locations captured from more than 1.2 million unique devices during a three-day period in 2017
(Source: New York Times 12/2018)



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33%/47% of **Android/iOS** apps shared **GPS coordinates** with third parties



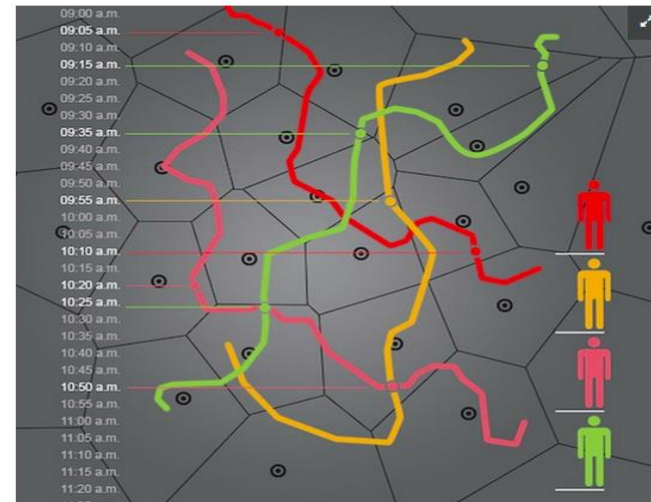
In about four months' of data reviewed by The Times, her location was recorded over 8,600 times — on average, once every 21 minutes.

TECHNOLOGY

LOCATION DATA CAN UNIQUELY IDENTIFY CELLPHONE USERS

A NEW STUDY DEMONSTRATES HOW EASY IT IS TO IDENTIFY PEOPLE FROM THE LOCATION TRACKING DATA ON THEIR CELLPHONES.

By Francie Diep March 27, 2013



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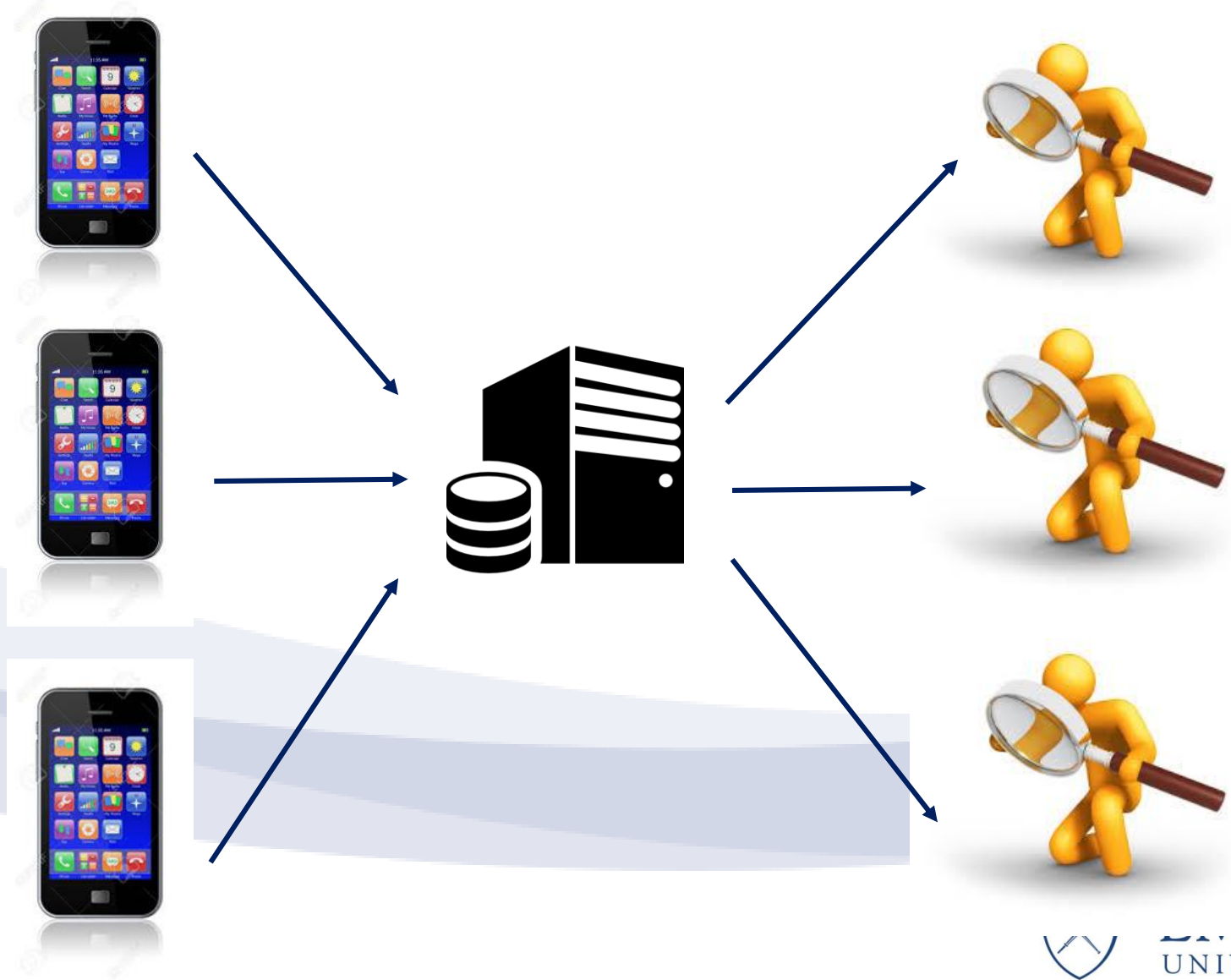
 Text Message Ukrainian Pro Cellphones Are Tracked

 New Tracking Pin Any Intern Location to Within Hundred Meters

 How Stores Track Shoppers

 Facebook Will Physically Store

The Mobile Data Economy



Enabling Data Analytics with **Centralized** Differential Privacy



Enable Mobile Apps and Analytics with **Local Differential Privacy**



Enabling Mobile Apps and Analytics with Local Differential Privacy

- Background
 - Local differential privacy
 - Geo-indistinguishability (local d-privacy)
- Extended privacy notions
 - Protecting dynamic locations (CCS15, VLDB17 demo)
 - Protecting spatiotemporal events (ICDE19)
- New mobile applications
 - Spatial crowdsourcing with geo-indistinguishability (ICDE18)
- New mechanisms
 - Supporting both analytics and mobile applications (CNS19)

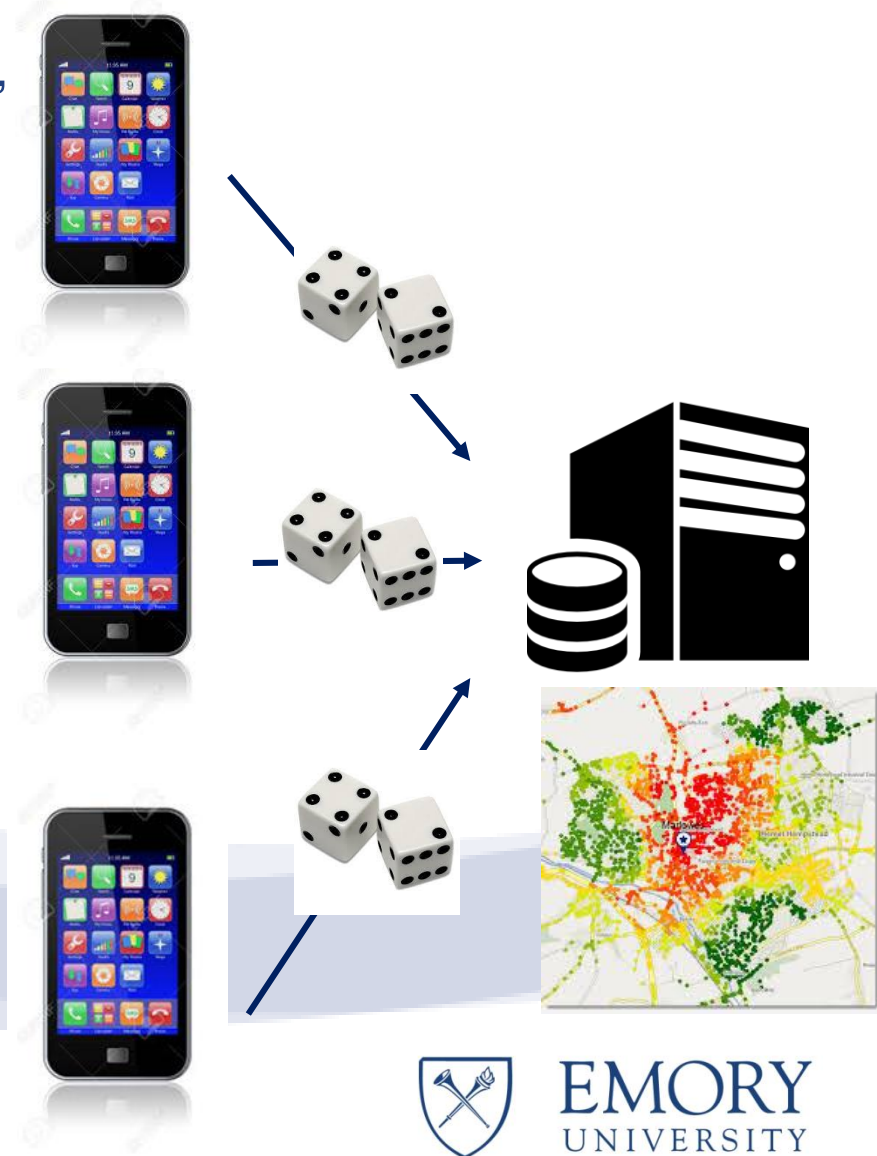


Local Differential Privacy

- Privacy definition
 - Any two locations produce “similar” distributions (bounded by ϵ)

$$\frac{Pr(\mathcal{A}(\mathbf{x}_1) = \mathbf{z}_t)}{Pr(\mathcal{A}(\mathbf{x}_2) = \mathbf{z}_t)} \leq e^\epsilon$$

- Mechanism
 - Randomized response (with encoding)
- Applications
 - Simple analytics (e.g. frequency estimation)
 - Google, Apple, Microsoft
- Limitations
 - Output not useful for mobile apps

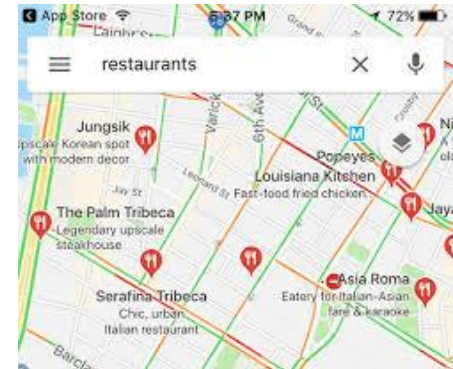
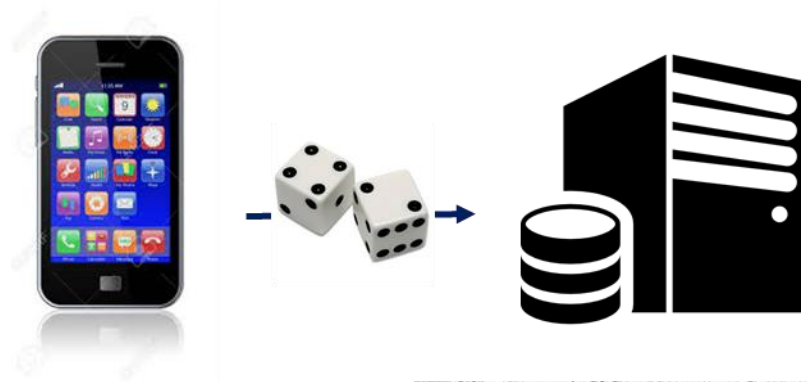


Geo-Indistinguishability (Local d-privacy)

- Privacy Definition
 - Any two locations **at distance at most r** produce “similar” distributions proportional to the distance (bounded by ϵr)

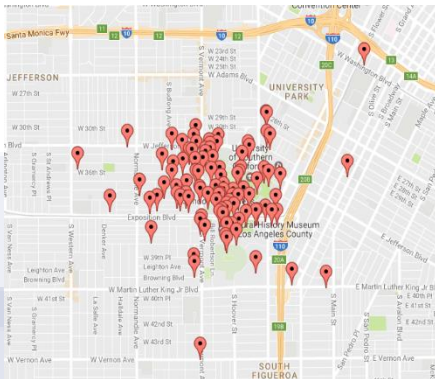
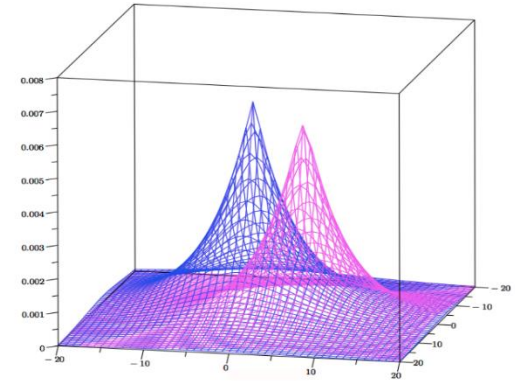
$$\frac{\Pr(\mathcal{A}(\mathbf{x}_1) = \mathbf{z}_t)}{\Pr(\mathcal{A}(\mathbf{x}_2) = \mathbf{z}_t)} \leq e^{\epsilon d(\mathbf{x}_1, \mathbf{x}_2)}$$

- Mechanism:
 - Planar Laplace mechanism
- Applications
 - Mobile apps/location sharing
- Limitations:
 - Temporal correlations of dynamic locations not considered
 - Not optimal for analytics

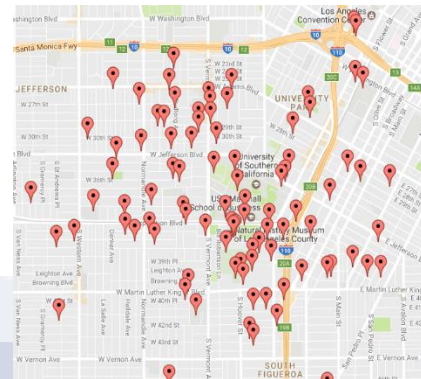


Geo-Indistinguishability: Planar Laplace Mechanism

- Generating random point z (from actual point $x \in X$) according to planar Laplace distribution



$\epsilon = \log(6)$
 $r = 1 \text{ km}$



Better privacy: $\epsilon = \log(2)$
 $r = 1 \text{ km}$



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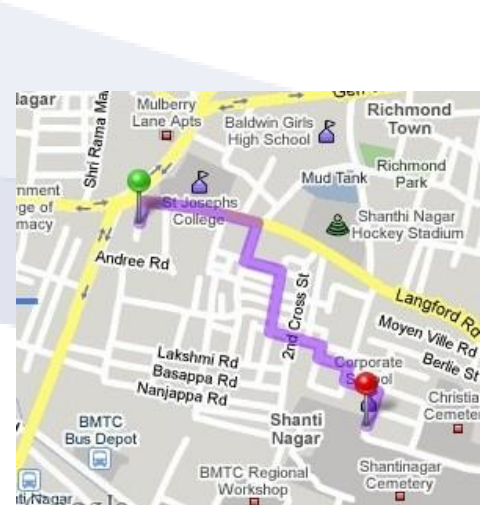
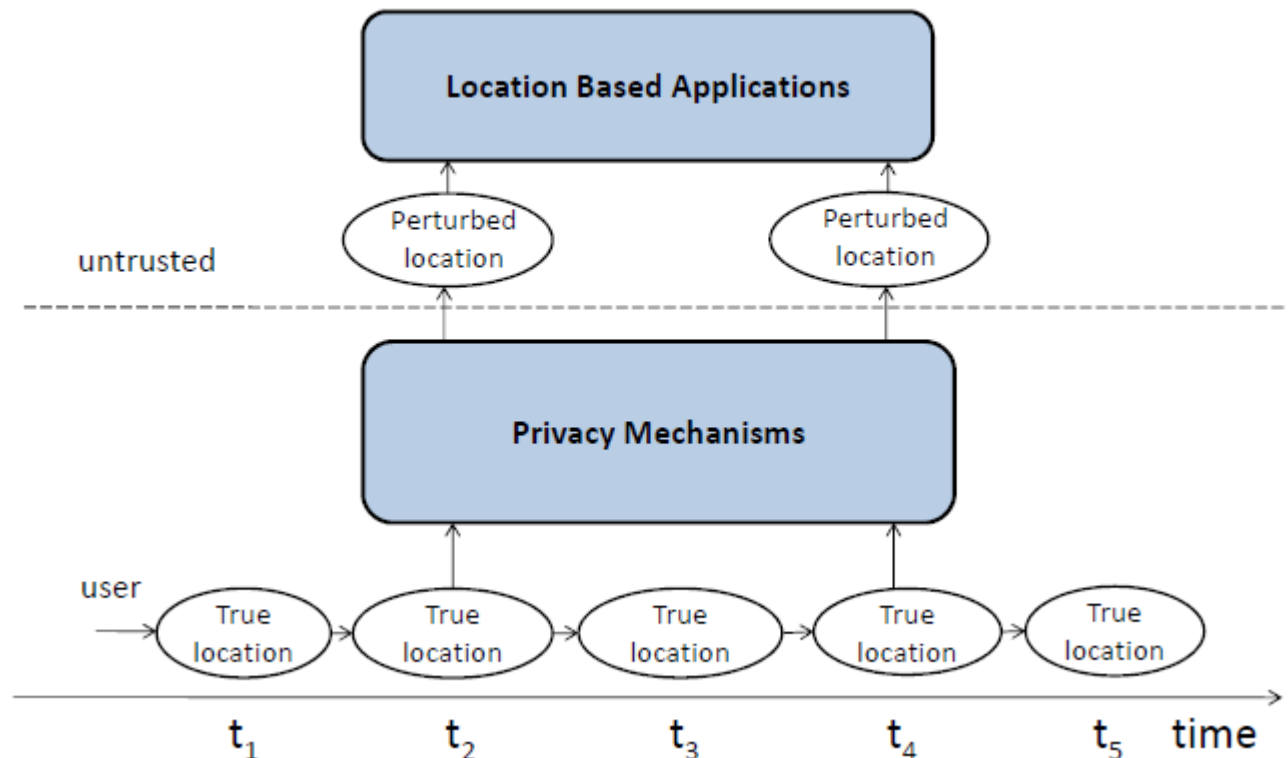
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Location Privacy: Temporal Correlations

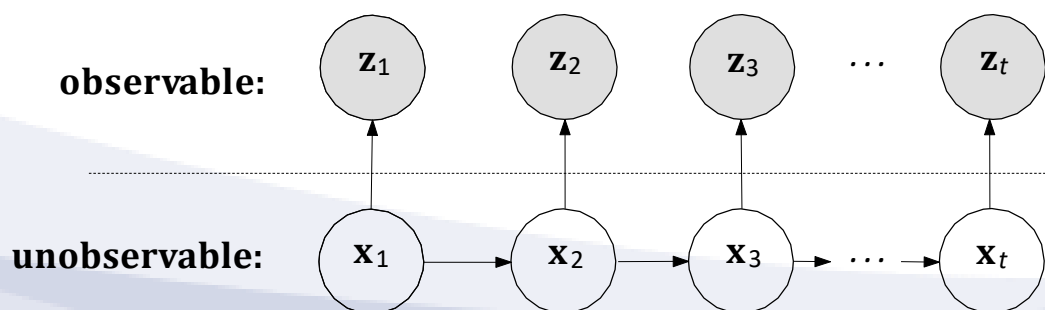
- Temporal correlations (adversary knowledge): moving patterns and previously released perturbed locations



Differential Privacy with δ -location set

- δ -location set differential privacy
 - Any two locations **in the probable location set** produce “similar” distributions proportional to the distance (bounded by ϵ)
 - Probable location set determined by hidden Markov Model

$$\frac{Pr(\mathcal{A}(\mathbf{x}_1) = \mathbf{z}_t)}{Pr(\mathcal{A}(\mathbf{x}_2) = \mathbf{z}_t)} \leq e^\epsilon$$



Y. Xiao, L. Xiong. Protecting Locations with Differential Privacy under Temporal Correlations. CCS 2015

Y. Xiao, L. Xiong, S. Zhang, Y. Cao. LocLok: Location Cloaking with Differential Privacy via Hidden Markov Model. VLDB demo, 2017



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Optimal perturbation mechanism

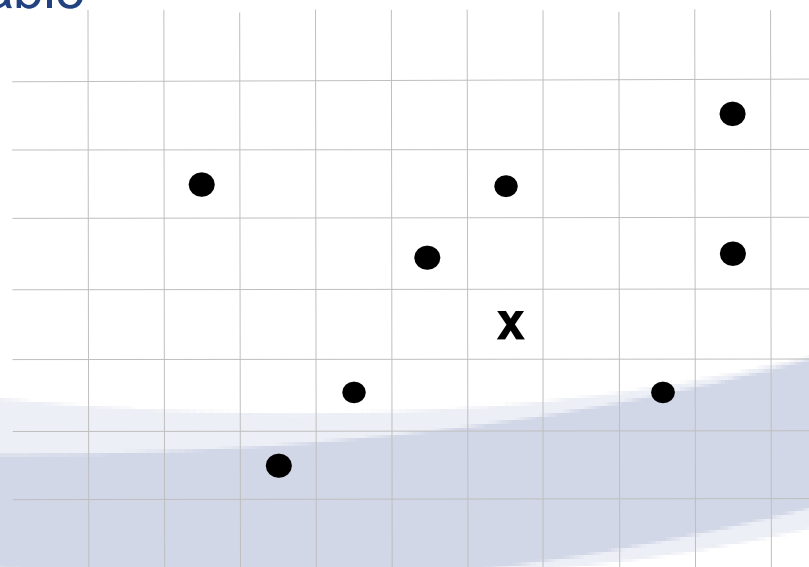
- Minimize **expected distance** between perturbed location \mathbf{z} and true location \mathbf{x}

$$\text{ERROR} = \sqrt{\mathbb{E} \|\mathbf{z} - \mathbf{x}^*\|_2^2}$$

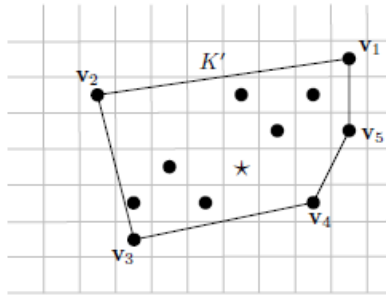
- While satisfying constraint of **differential privacy** – any pair of locations \mathbf{x}_1 and \mathbf{x}_2 are indistinguishable

$$\frac{\Pr(\mathcal{A}(\mathbf{x}_1) = \mathbf{z}_t)}{\Pr(\mathcal{A}(\mathbf{x}_2) = \mathbf{z}_t)} \leq e^\epsilon$$

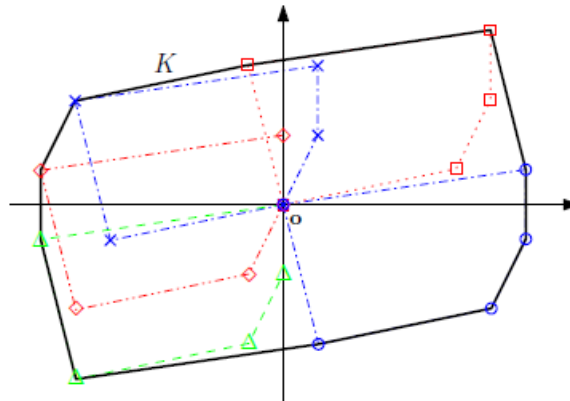
- Exponential mechanism and Laplace mechanism are not optimal



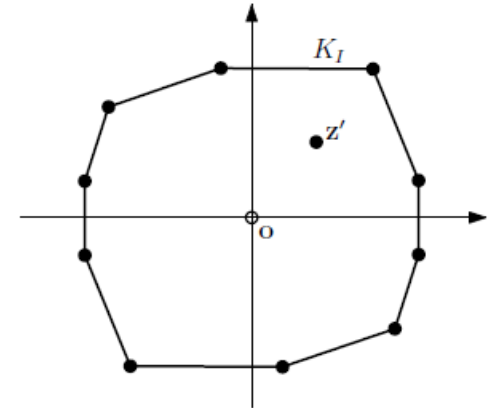
Planar Isotropic Mechanism



(a)



(b)

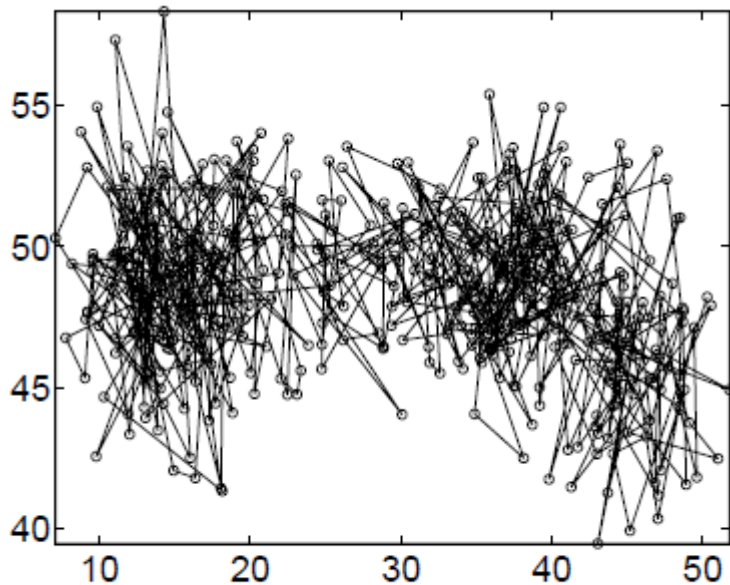
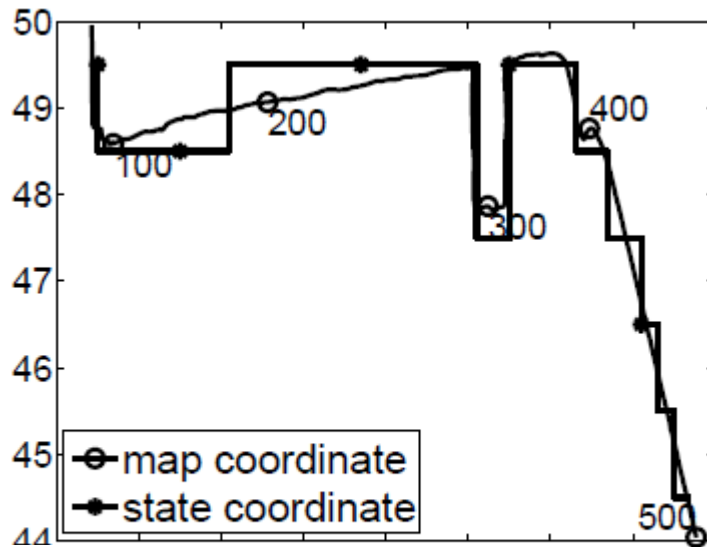


(c)

- Based on sensitivity hull K of δ -location set which determines the lower bound error
- An improved K -norm mechanism based on Isotropic transformation
- Achieves optimality while achieving differential privacy

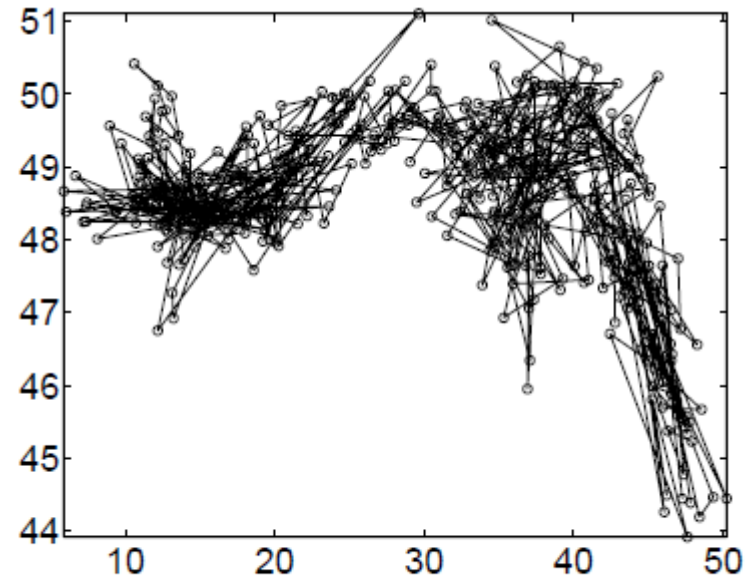


Results: Perturbed Trace Illustration



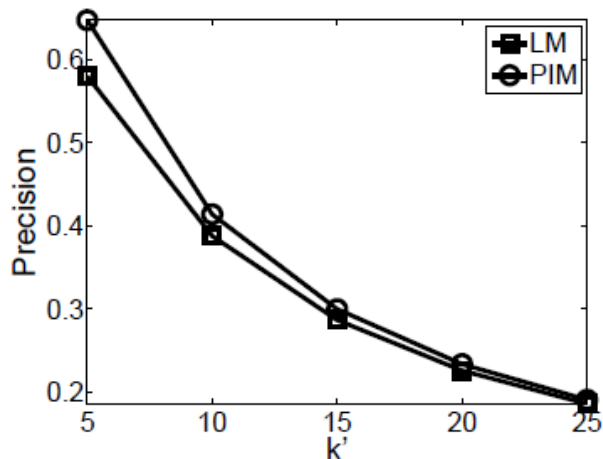
(c) LM released trace

(a) Tru

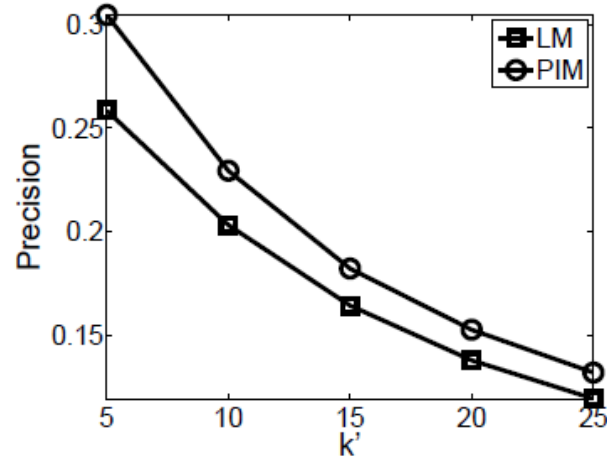


(e) PIM released trace

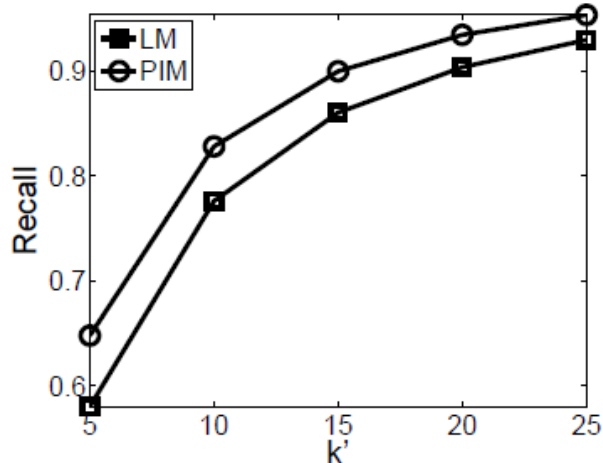
Results: k-Nearest Neighbor Queries



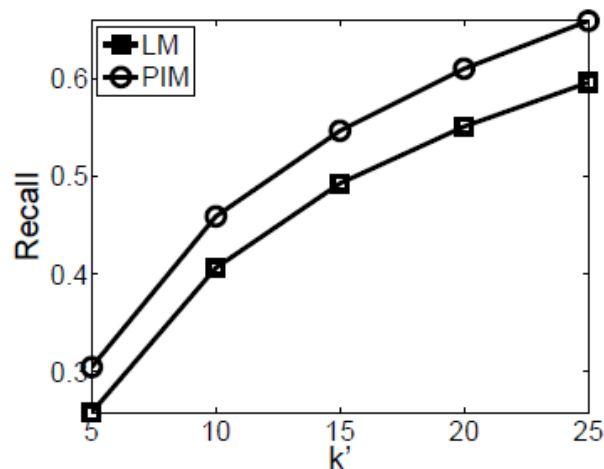
(c) Precision (Geolife)



(d) Precision (Gowalla)



(e) Recall (Geolife)



(f) Recall (Gowalla)

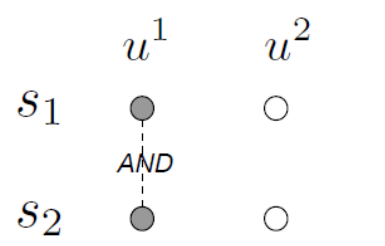
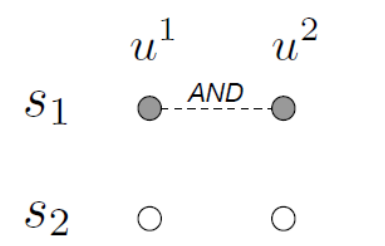
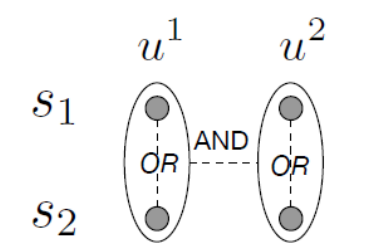
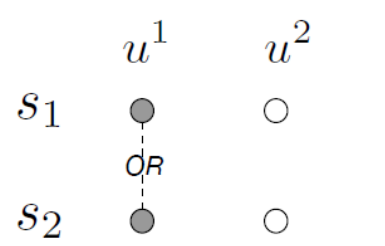
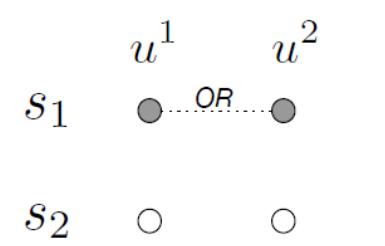
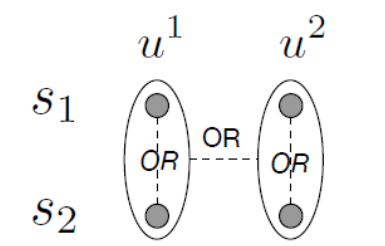
From Location Privacy to Spatiotemporal Privacy

- Location privacy mechanisms protect **location at a time point**
- May not protect **spatiotemporal activities?**
 - Staying in hospital for 2 hours
 - From home to office every morning
- Need formal notions and mechanisms



Spatiotemporal events

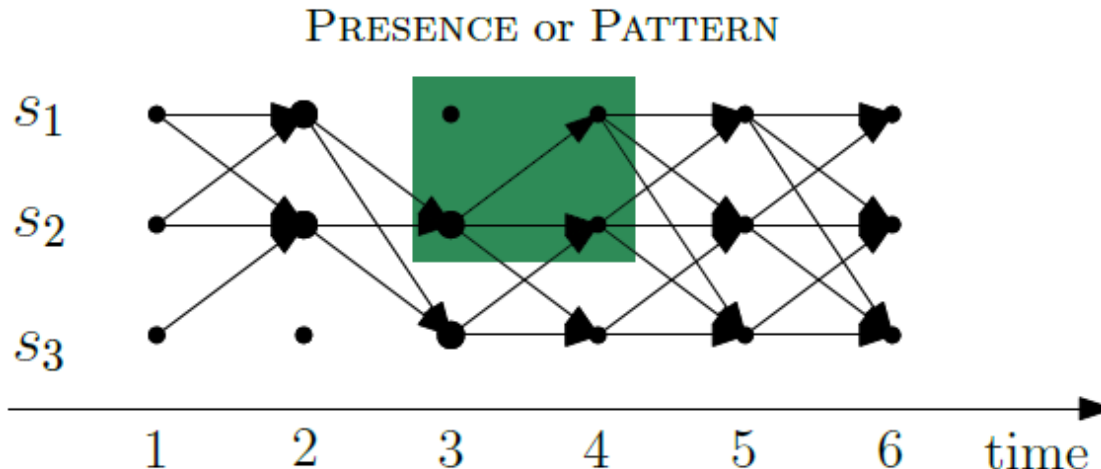
- Boolean expression for spatiotemporal event
- Location at a time point ($u^t = s_i$)

Spatial dimension	Temporal dimension	Spatial and Temporal
 <p>(a) $(u^1 = s_1) \wedge (u^1 = s_2)$</p>	 <p>(c) $(u^1 = s_1) \wedge (u^2 = s_1)$</p>	 <p>(e) $((u^1 = s_1) \vee (u^1 = s_2)) \wedge ((u^2 = s_1) \vee (u^2 = s_2))$</p>
 <p>(b) $(u^1 = s_1) \vee (u^1 = s_2)$</p>	 <p>(d) $(u^1 = s_1) \vee (u^2 = s_1)$</p>	 <p>(f) $((u^1 = s_1) \vee (u^1 = s_2)) \vee ((u^2 = s_1) \vee (u^2 = s_2))$</p>

From Location Privacy to Spatiotemporal Event Privacy

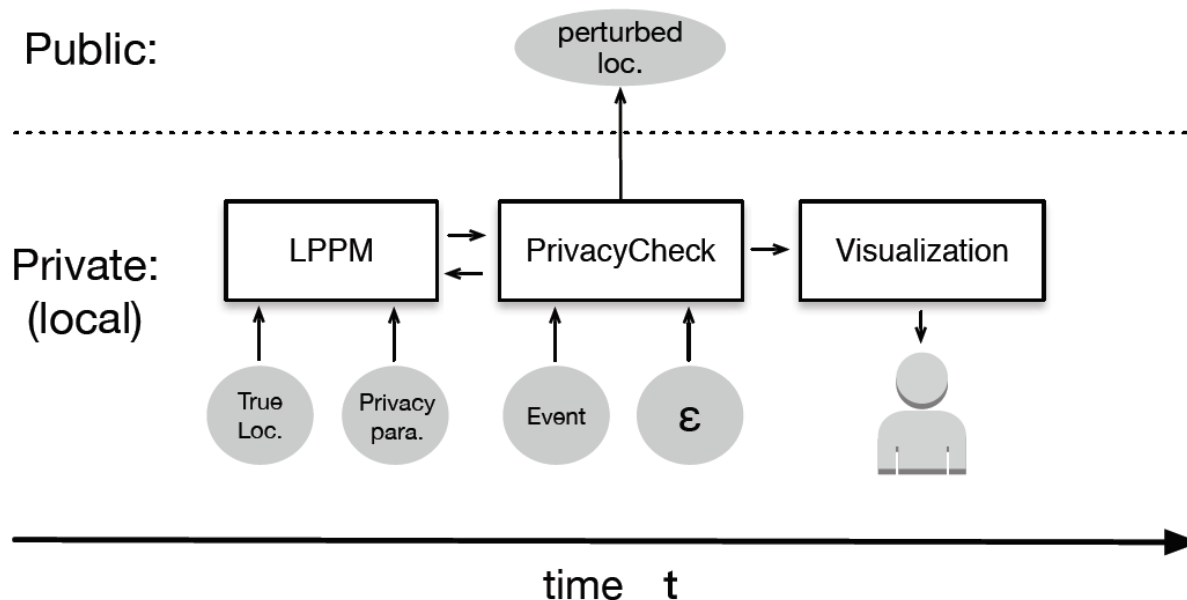
- Location privacy
 - **Two locations** produce “similar” distributions/observations
- Spatiotemporal event privacy
 - **A true event and a negative event** produce “similar” location traces

$$Pr(o_1, o_2, \dots, o_t | \text{EVENT}) \leq e^\epsilon Pr(o_1, o_2, \dots, o_t | \neg \text{EVENT})$$



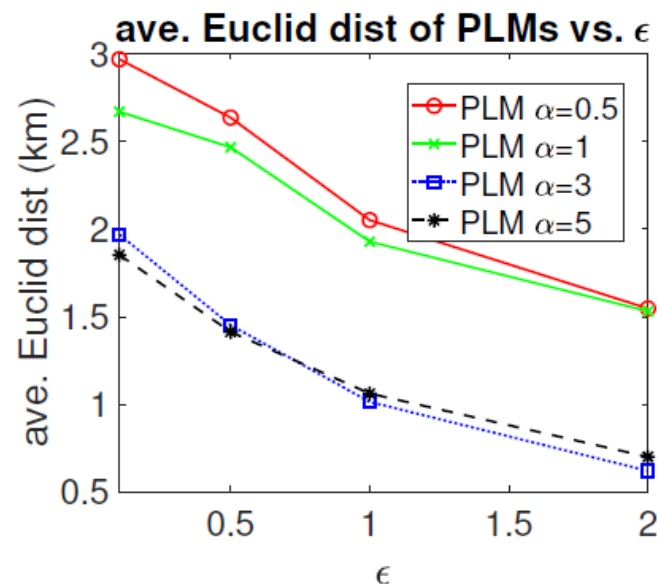
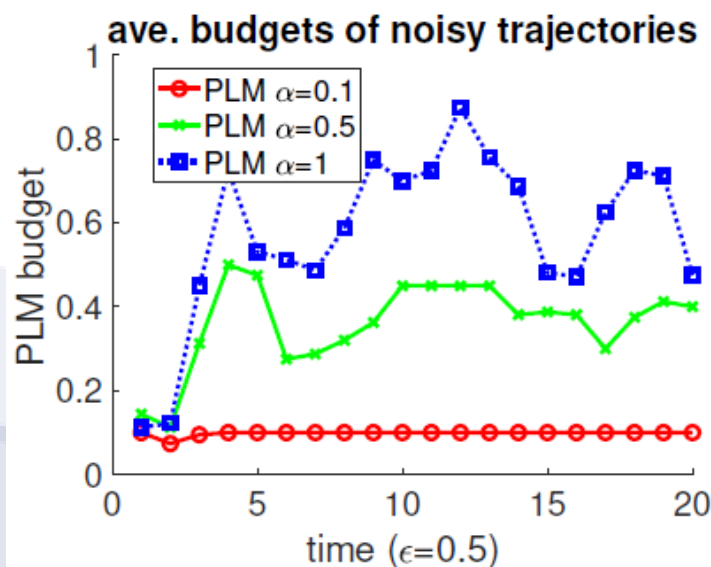
Spatiotemporal Privacy Framework

- **LPPM**: Existing location privacy mechanism, e.g. Planar Laplace Mechanism for geo-indistinguishability
- **PrivacyCheck**: check spatiotemporal event privacy and calibrate privacy budget



Results

- Strong LPPM may satisfy spatiotemporal privacy already
- Weak LPPM need to reduce privacy budget significantly (less utility) to achieve same level of spatiotemporal privacy
- Stronger spatiotemporal privacy, less utility of the locations

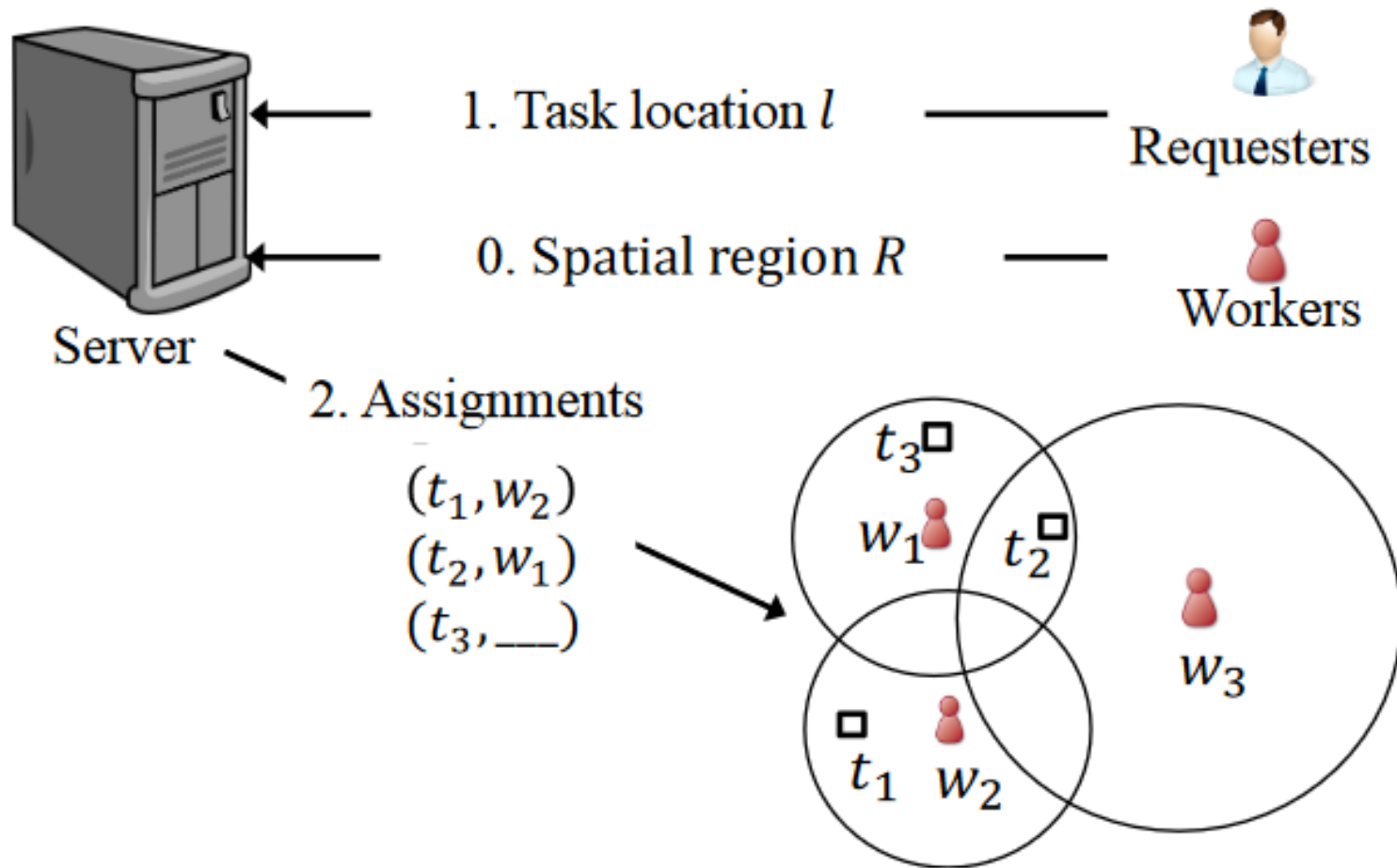


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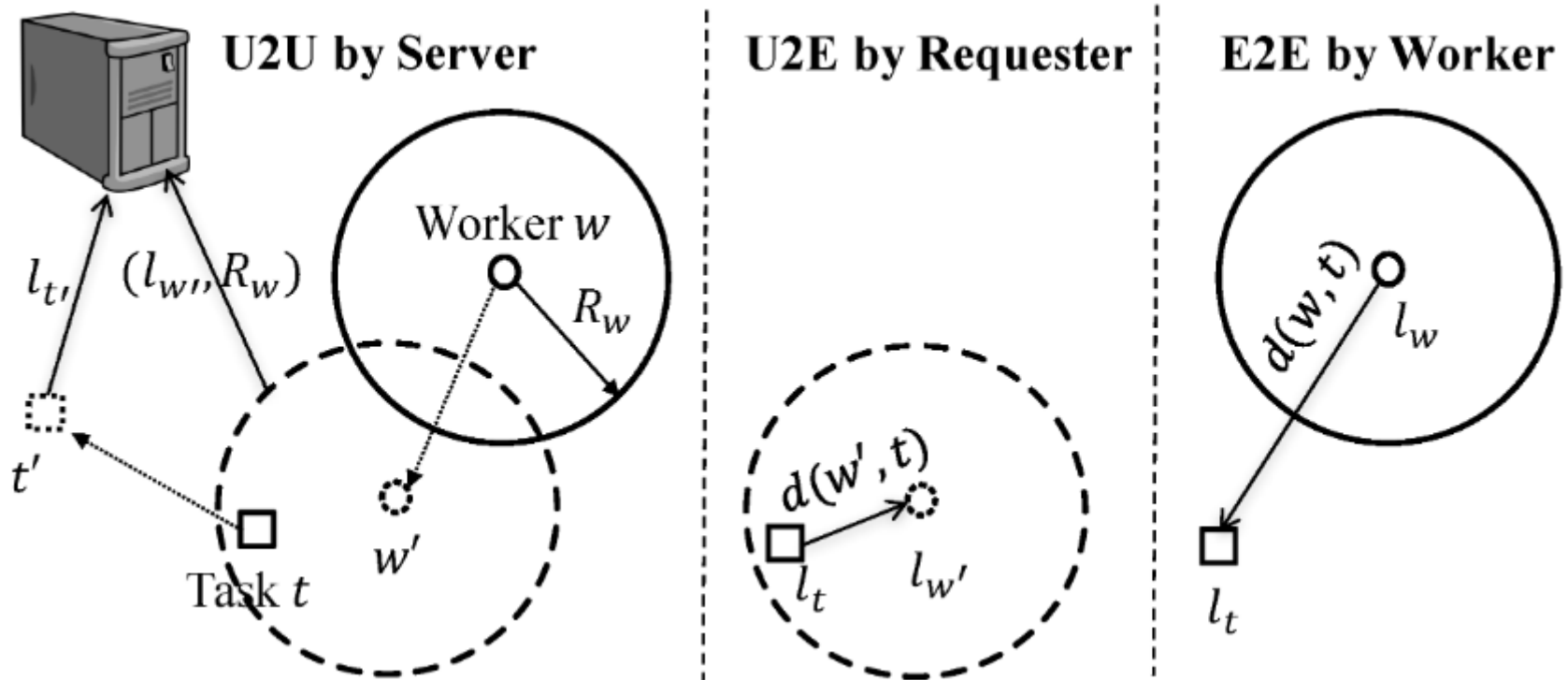


ONLINE TASK ASSIGNMENT IN SPATIAL CROWDSOURCING



Privacy preserving online task assignment in spatial crowdsourcing

- Both requester and worker locations are perturbed using **geo-indistinguishability**
- Three-stage framework for task assignment using uncertain locations



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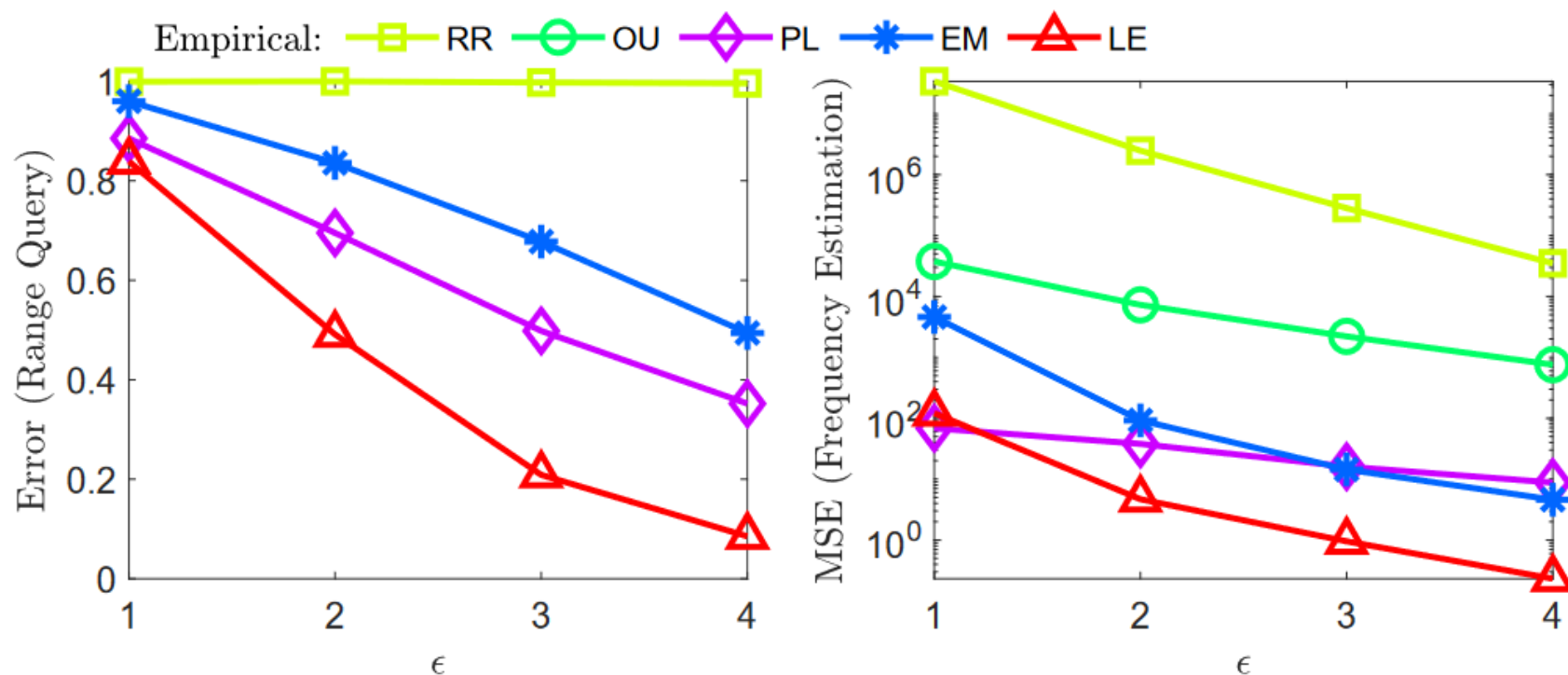
Supporting both range queries and frequency estimation

- Existing
 - Local differential privacy with randomized response – frequency estimation
 - Geo-indistinguishability (local d-privacy) with planar Laplace mechanism – range queries
- Goal
 - Optimize for both frequency estimation and range queries while ensuring local d-privacy
- Basic idea
 - Assign different perturbation probabilities for different input/output pairs in a way related to the distance



Results: Comparison

Gowalla dataset



RR: Randomized Response
OU: Optimized with Unary Encoding
PL: Planar Laplace mechanism
EM: Exponential mechanism
LE: Linear equation mechanism



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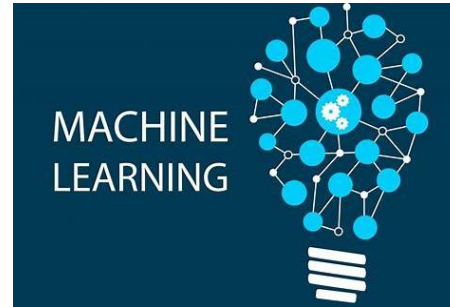
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- Open challenges
 - Privacy/utility tradeoff
 - User empowerment



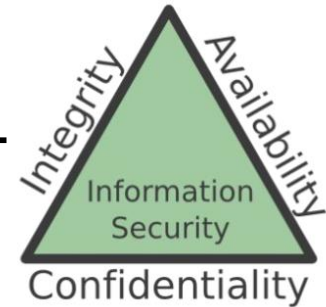
Assured Information Management and Sharing (AIMS)



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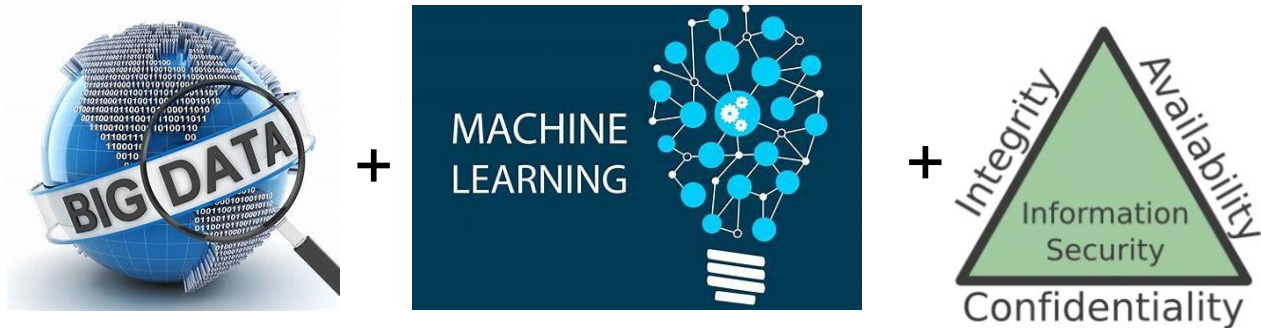


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Assured Information Management and Sharing (AIMS)



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