

# Community Search over Big Graphs: Models, Algorithms, and Opportunities

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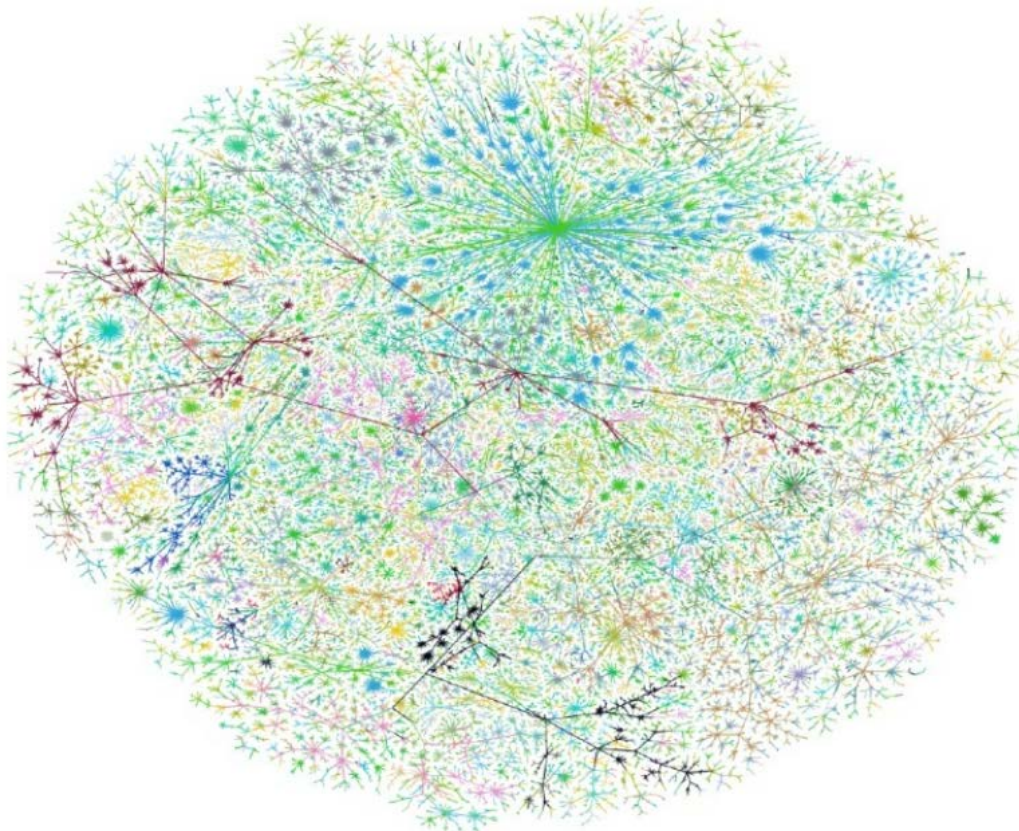


# Tutorial Outline

- Introduction, Motivations, and Challenges
- Networks & Community Detection
- Community Search (4 Parts)
  - Densely-connected community search
  - Attributed community search
  - Social circle discovery
  - Querying geo-social groups
- Future Work & Open Problems

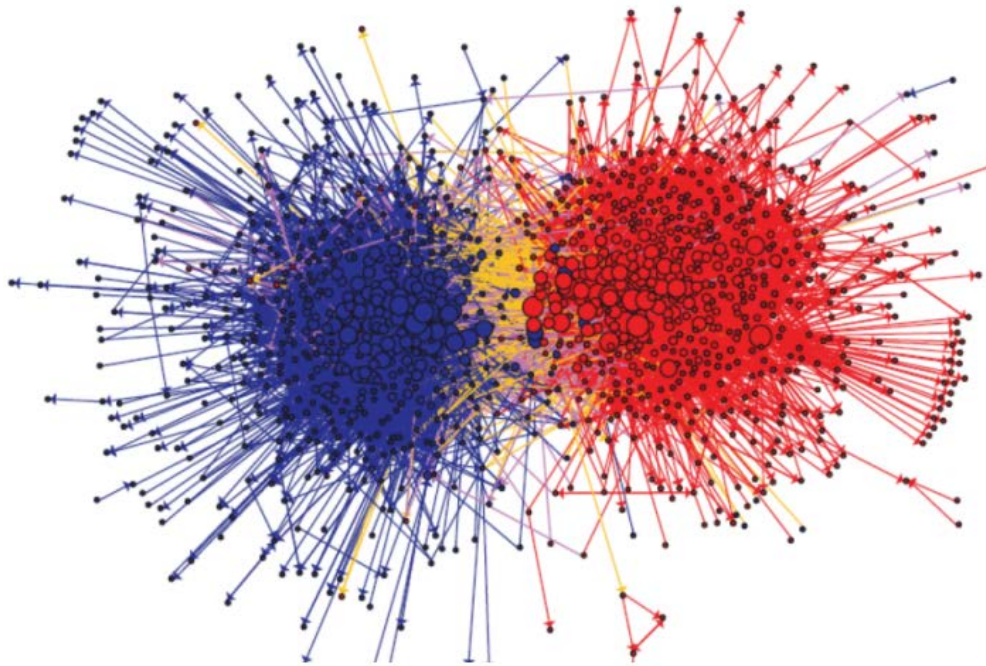
# Networks

- **Networks** are everywhere (e.g. chemistry, biology, social networks, the Web, etc.)



# Communities

- **Communities** naturally exist in **networks**.



**Blogosphere**

# Community Structure

- **Community structure:** Nodes with a shared latent property, densely inter-connected .
- Many reasons for communities to be formed:

Social Networks



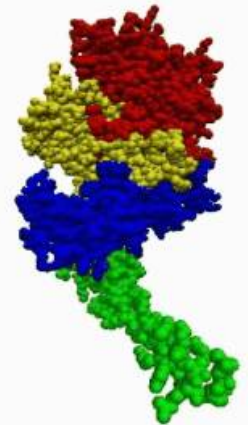
Citation Networks



World Wide Web

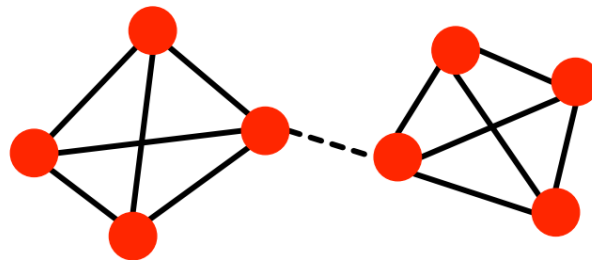


Biological Networks



# Basis of Community Formation

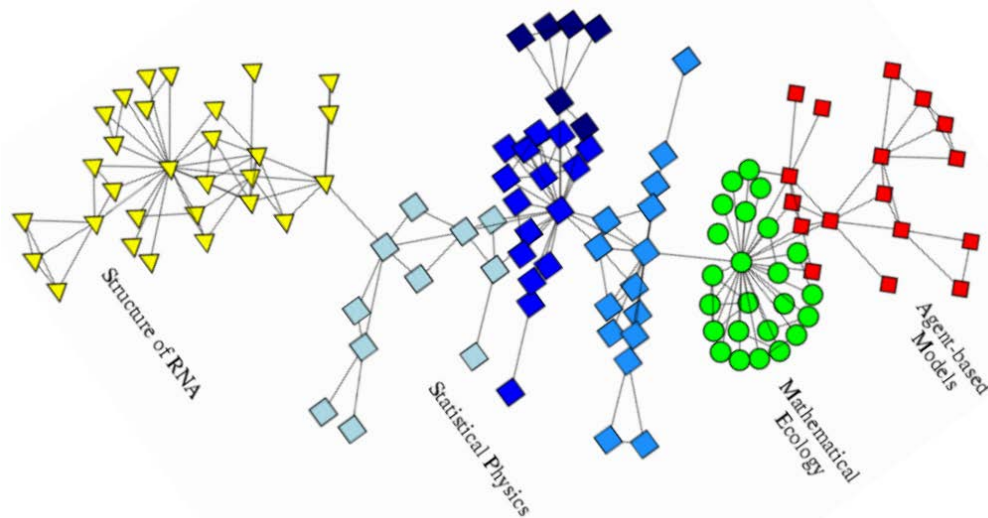
- **The strength of weak ties** [Mark Granovetter, 1973] and **the models of small-world** [Strogatz and Watts, Nature'98] both suggest
  - **Strong ties** are well embedded in the network
  - **Weak ties** span long ranges



- **Given a network, how do we find all communities?**

# Community Detection

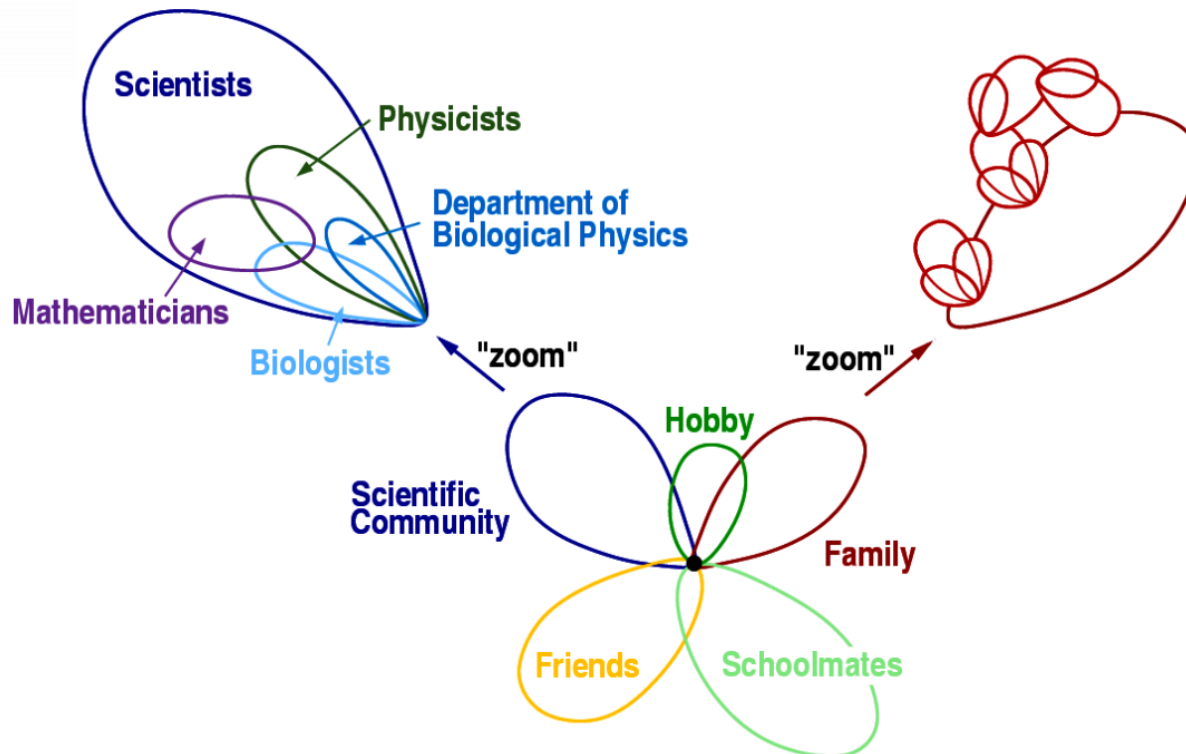
- **Q: Given a network, how do we find all communities?**
- **A: Find weak ties and identify communities**
  - **Betweenness centrality** [Girvan and Newman, PNAS'02],
  - **Modularity** [Newman, PNAS'06]
  - **Graph partitioning methods** [Karypis and Kumar, SISC'08]



SFI collaboration network [Newman]

# Overlapping Communities

- **Communities** defined by **different nodes** in a network may be quite different.





# Community Search

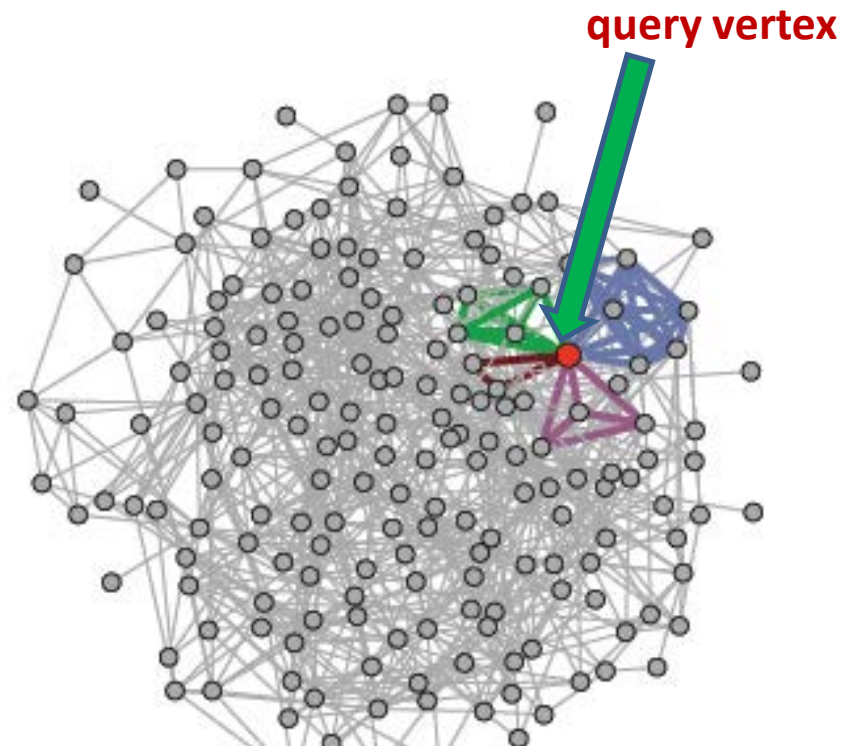
- **Problem:** Given a set of **query nodes**, find densely connected communities containing them.

- State-of-the-art research focus:

**Simple** and **static** graphs →

**Evolving, attributed,** and

**location-based** big graphs



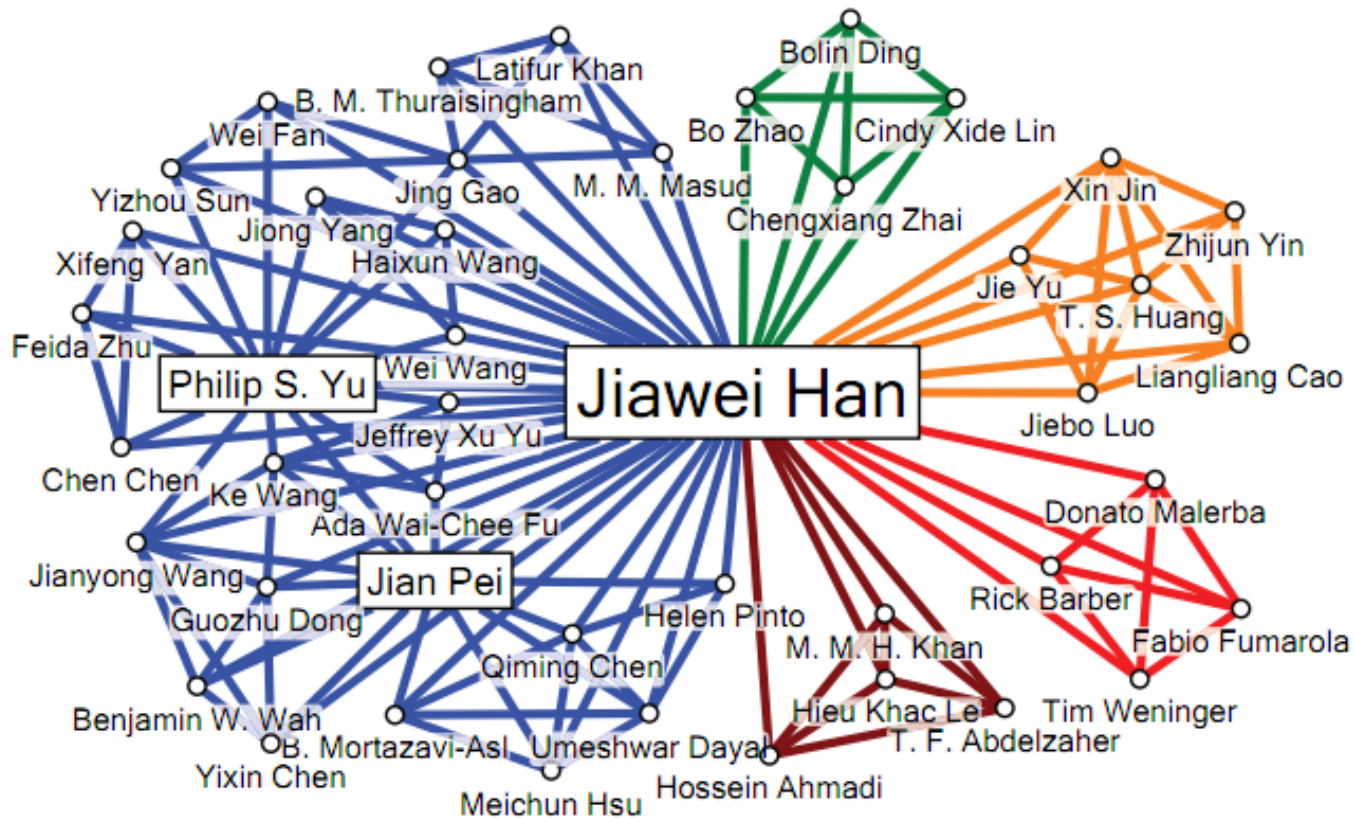
# Community Detection v.s. Community Search

- **Community detection: identify all communities.**
  - fundamental & widely studied
  - global computation (expensive)
  - static graphs (hard to handle evolving graphs)
- **Community search: find query-dependent communities**
  - useful & less studied
  - user-centered & personalized search
  - dynamic graphs

# Applications

- **Social circle discovery**
- **Planning a cocktail party/conference/workshop**
- **Infectious disease control**
- **Tag recommendation**
- **Protein complex identification**

# Community Search



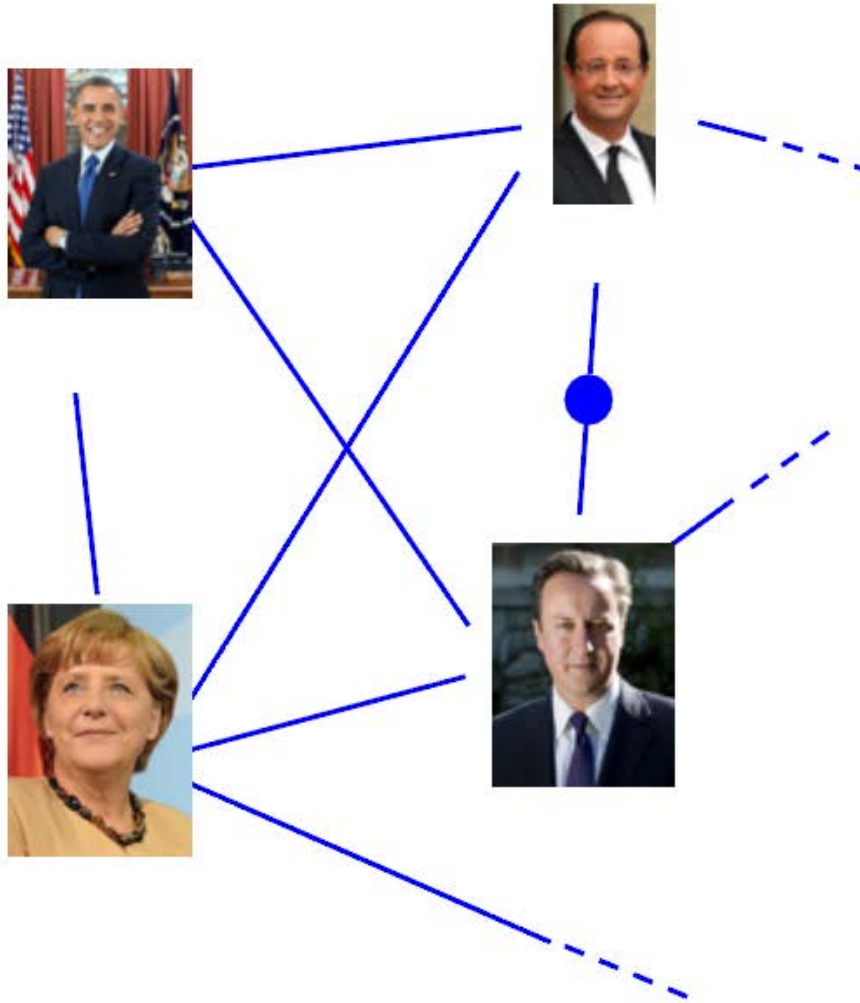
5 communities containing “Jiawei Han” in DBLP collaboration network

# Planning a cocktail party

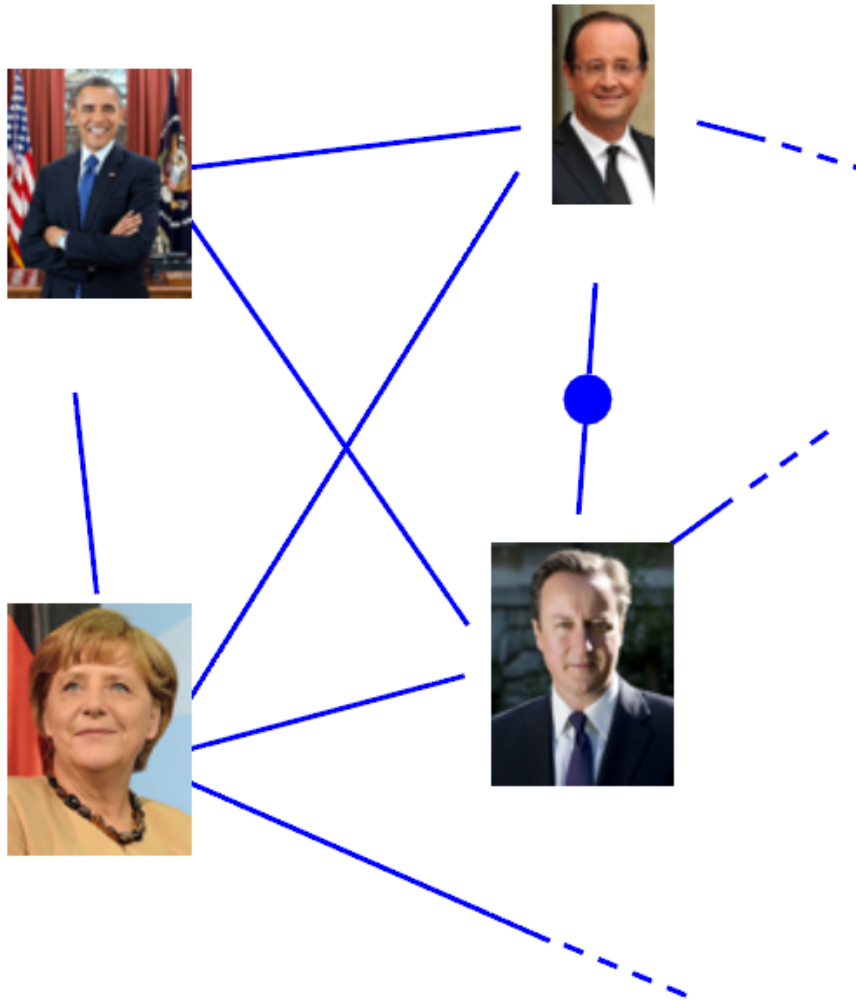
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# Planning a cocktail party



# Planning a cocktail party



Recipe for a successful party:

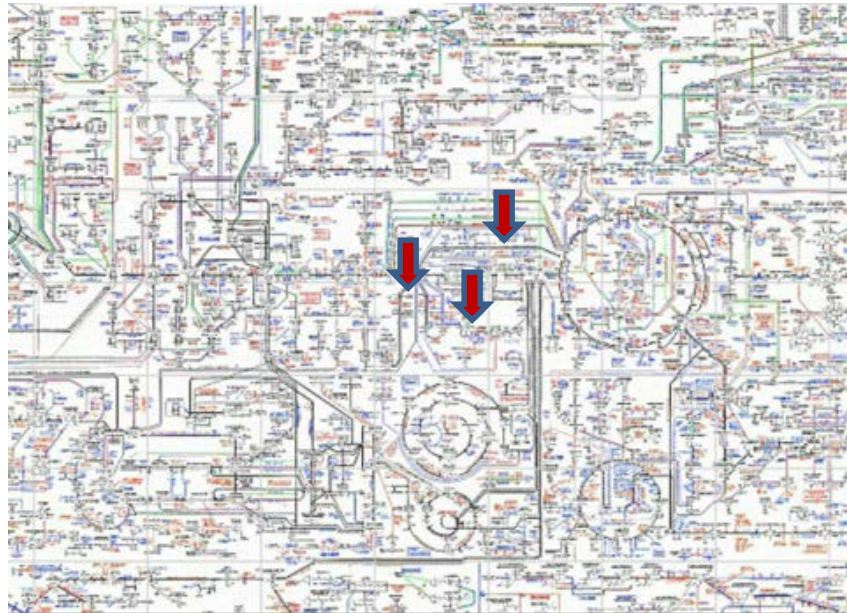


- Participants should be “close” to the organizers (e.g., a friend of a friend).
- Everybody should know sufficiently many in the party (on an average?).
- The graph should be connected.
- The number of participants should not be too small but...
- ...not too large either!!!
- ....
- social distance not too large.

Not an easy task...

# Protein Complex Identification

- Given: a protein-protein interaction network
- A set of proteins that regulate a gene that a biologist wishes to study.
- What other proteins should she study? those contained in a compact dense subgraph containing the given proteins.





# Challenges

- Complexity of underlying community models
- Responsiveness requirements of query processing
- Dynamic network structures
- Massive volume of big graphs

# Related Work

- **Community Detection** (Finding all communities in the entire network)
  - non-overlapping community detection [*Girvan and Newman, PNAS'02*]
  - overlapping community detection [*Ahn et al, Nature'10*]
- **Community Search** (Finding communities containing given query nodes)

Different community models are proposed for various types of networks and query processing techniques.

  - **Structural Networks** ---> Densely-connected community search
  - **Attributed Graphs** ---> Attributed community search
  - **Ego-networks** ---> Social circle discovery
  - **Location-based Social Networks** ---> Querying geo-social groups

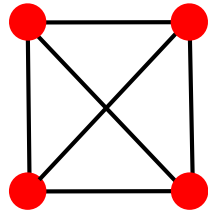
# Part 1: Densely-connected Community Search

- In the simplest way, a graph represents a structure of interactions within a group of vertices.
- Task: finding **densely-connected communities containing query nodes**.
  - **Quasi-clique model** [Cui et al. SIGMOD'13]
  - **Query-biased densest subgraph model** [Wu et al. PVLDB'15]
  - **K-core model** [Sozio & Gionis KDD'10, Cui et al. SIGMOD'14, Li et al. PVLDB'15, Narbieri et al. DMKD'15]
  - **K-truss model** [Huang et al. SIGMOD'14, Huang et al. PVLDB'16]

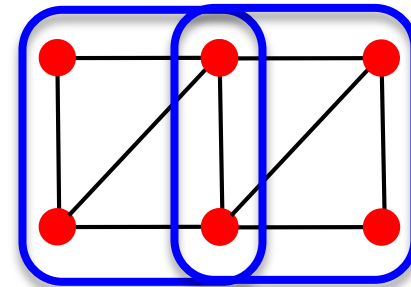
# Quasi-Clique based Model

- **$\alpha$ -adjacency- $\gamma$ -quasi- $k$ -clique community model**
  - **$\gamma$ -quasi- $k$ -clique**: a  $k$ -node graph with at least  $\lfloor \gamma k(k-1)/2 \rfloor$  edges.
  - **$\alpha$ -adjacency- $\gamma$ -quasi- $k$ -clique**: overlap  $\alpha$  vertices, where  $\alpha \leq k-1$ .

**$k$ -clique**: a complete graph of  $k$  nodes with  $k(k-1)/2$  edges.



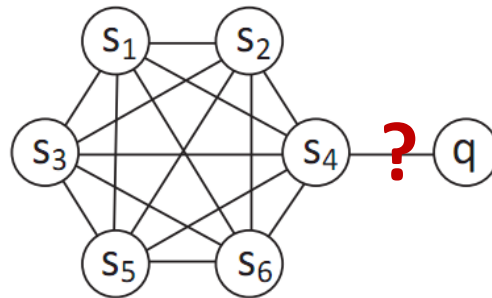
$\gamma$ -quasi- $k$ -cliques  
( $\gamma=1, k=4$ )



$\alpha$ -adjacency- $\gamma$ -quasi- $k$ -cliques  
( $\alpha=2, \gamma=0.8, k=4$ )

# Quasi-Clique based Model

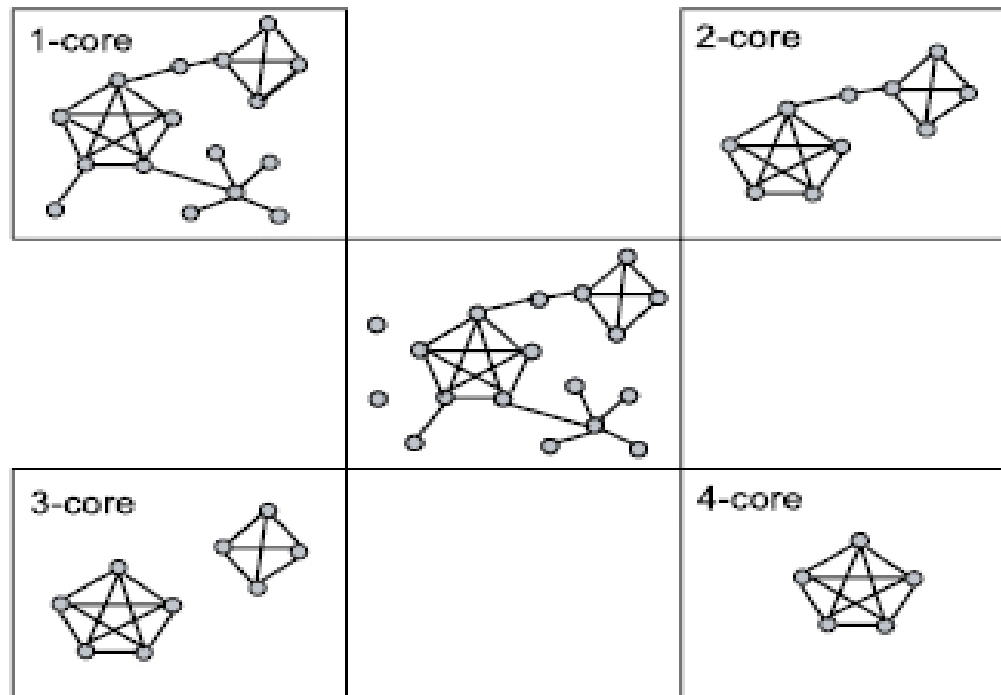
- **Problem:** Given a query vertex  $q$  in graph, the problem is to find all  $\alpha$ -adjacency- $\gamma$ -quasi- $k$ -clique containing  $q$ .



A 0.8-quasi-7-clique containing  $q$

# K-Core

- **K-core**: every vertex has degree *at least*  $k$  in this subgraph.

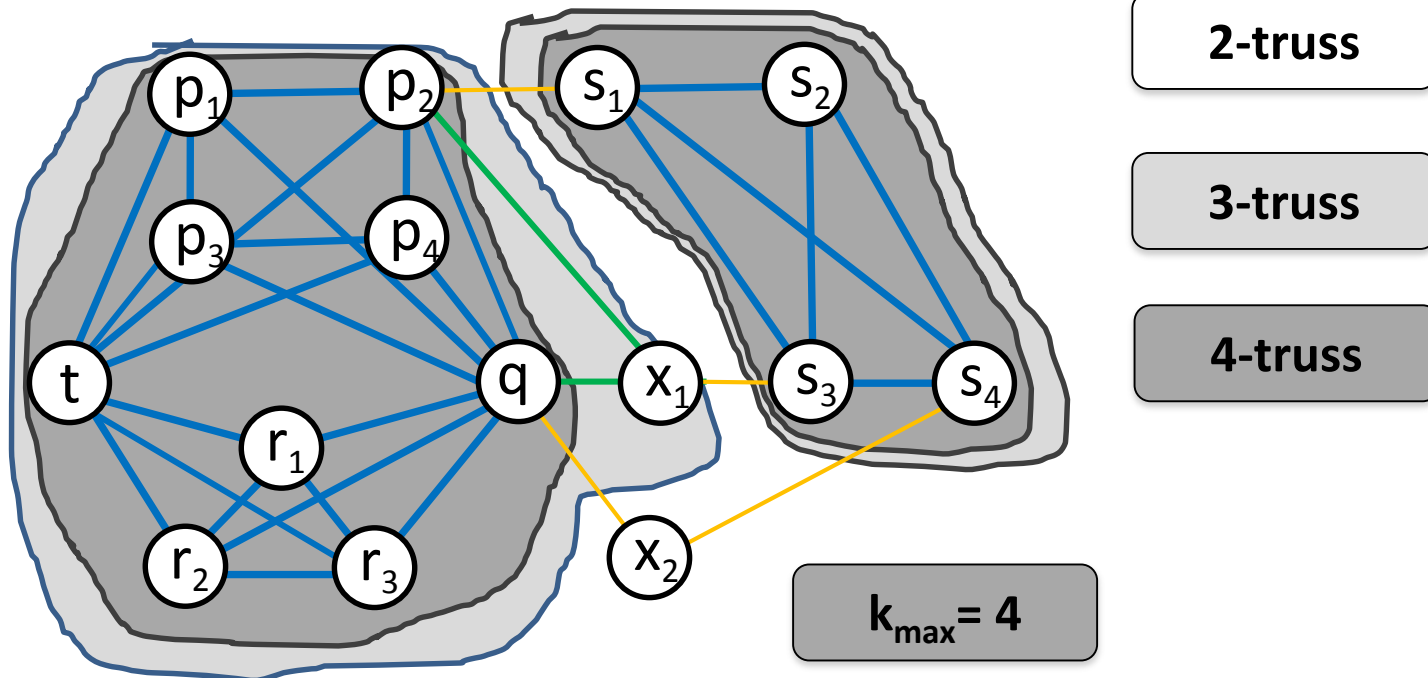


# K-Core based Model

- Input:  
a graph  $G$  & a set of query nodes  $Q$
- Output: a connected subgraph  $H$  containing  $Q$  such that
  - (1) Query distance  $D_Q(H) \leq \text{distance constraint}$ .
  - (2)  $|V(H)| \leq \text{size constraint}$ .
  - (3)  $H$  is a  $k$ -core with the largest  $k$  by satisfying (1) and (2).
- Other  $k$ -core based community models:
  - Local search algorithm [Cui et al. SIGMOD'14]
  - Minimum-size Community [Narbieri et al. DMKD'15]
  - Influential Community [Li et al. PVLDB'15]

# K-Truss

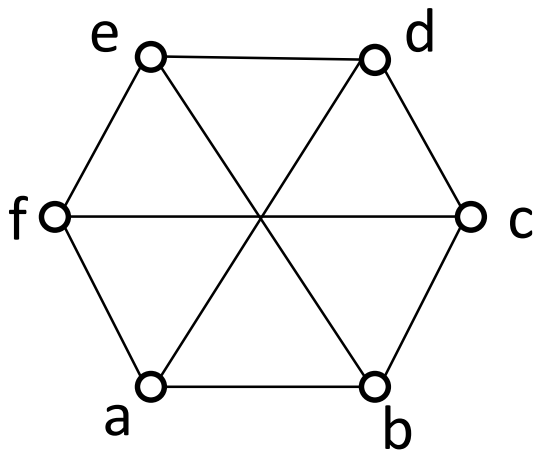
- **Triangle:** fundamental **building blocks** of networks
- **k-truss** of graph G: **every edge** in H is contained in **at least (k-2) triangles** within H.



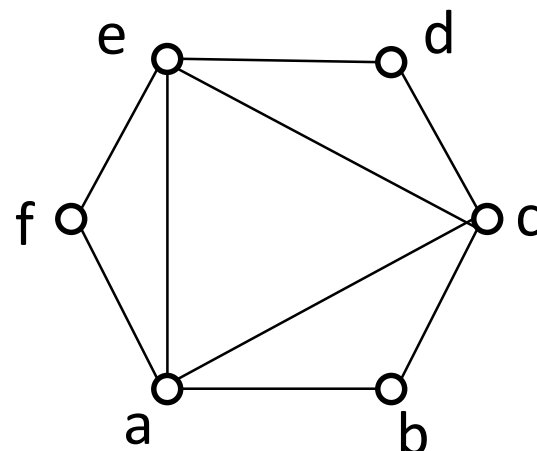


# K-Core V.S. K-Truss

- **K-core:** any pair of vertices within an edge may have **no common neighbors**.
- **K-truss:** any pair of vertices within an edge must have  **$k-2$  common neighbors**.



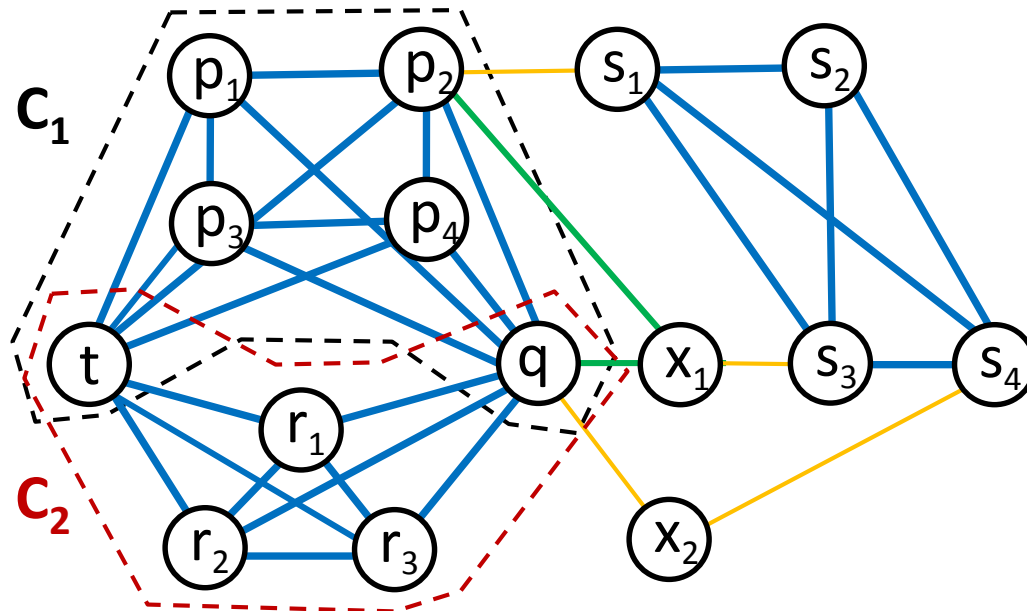
**3-core**



**3-truss**

# K-truss Community Model

- **A k-truss community** satisfies:
  - (1) **K-truss**: each edge within *at least (k-2) triangles*
  - (2) **Edge Connectivity**: all pairs of edges connected by triangles
  - (3) **Maximal Subgraph**



Two 4-truss communities for  $q$

# Problem Formulation

- **Problem:** Given a graph  $G(V, E)$ , a query vertex  $q$  and an integer  $k \geq 3$ , find all  $k$ -truss communities containing  $q$ .

# Index Based Query Processing Algorithm Framework

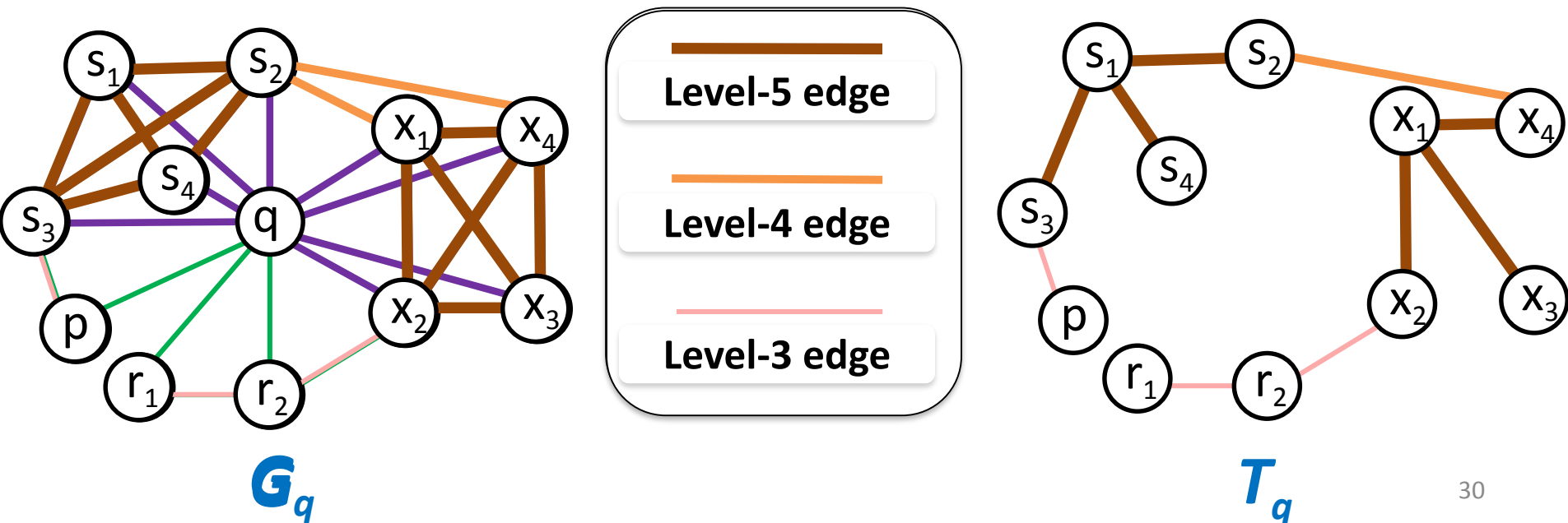
- Several different **index structures** are designed for the efficient search of **k-core** and **k-truss** based communities.
- We take the **k-truss community model** as an example.

# Index Based Query Processing Algorithm Framework

- **Index Construction (offline)**
  - They design a novel and compact **tree-shaped structure** called **TCP-index**.
- **Query Processing (online)**
  - Based on **TCP-index**, k-truss community search can be done in **optimal time** complexity.

# TCP-Index Construction

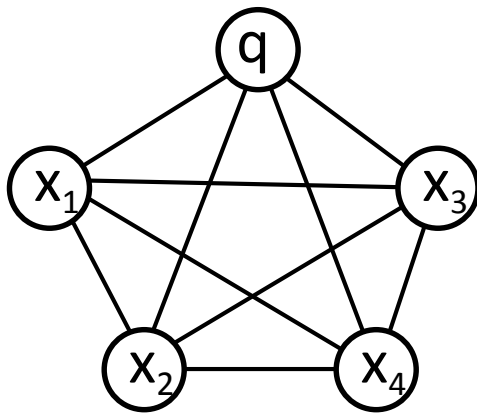
- TCP-Index for vertex  $x$  is a **tree structure as  $T_x$** .
  - $T_x$  is a **maximum spanning forest**.
  - Build  $T_x$  with weighted edges **level by level**.
  - $O(m)$  linear disk space,  $O(|Ans|)$  optimal query time.



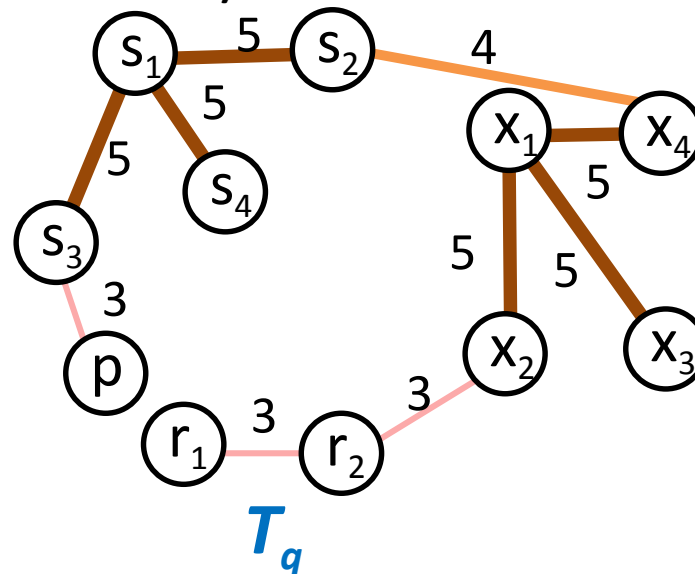
# Query Processing using TCP-Index

- Rationale:** If  $y, z$  are connected via a series of edges with weight  $\geq k$  in  $T_x$ , then  $y, z$  are in the same  $k$ -truss community; We use  $V_k(x, y)$  to denote all such vertices  $z$ .
- For example, querying **5-truss communities containing  $q$** .

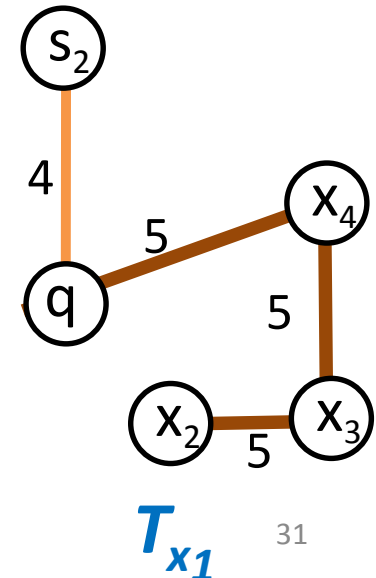
Each edge is accessed only 2 times. **Constant!!!**  
 (First time in black; Second time in red.)



A complete graph  $G = (V, E)$  with  $V = \{q, x_1, x_2, x_3, x_4\}$



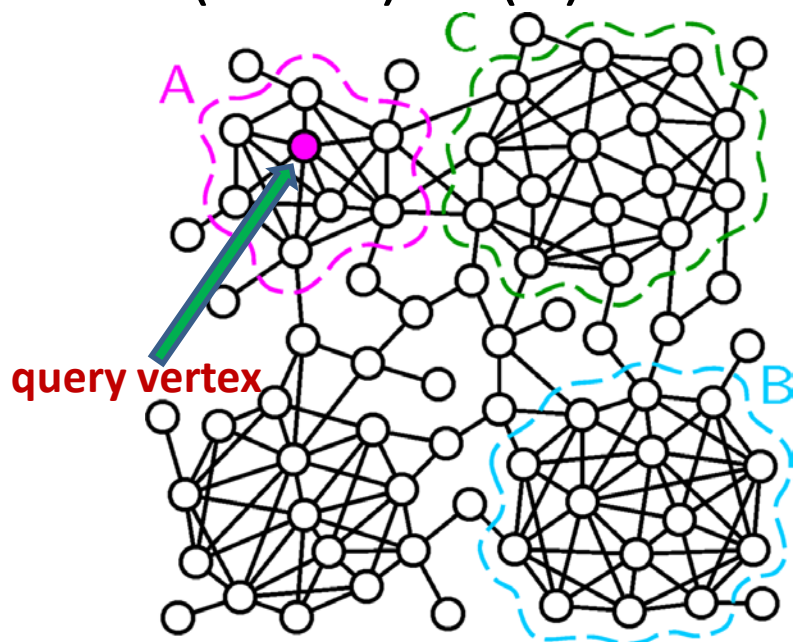
$T_q$



$T_{x_1}$

# Motivation: Free Rider Effect

- **Free Rider Effect:** far away and irrelevant nodes are included into communities.
- **Classic density:**  $f(S) = |E(S)| / |S|$ ,  $E(S) = E \cap S^2$
- $f(A \cup B) > f(A)$ .



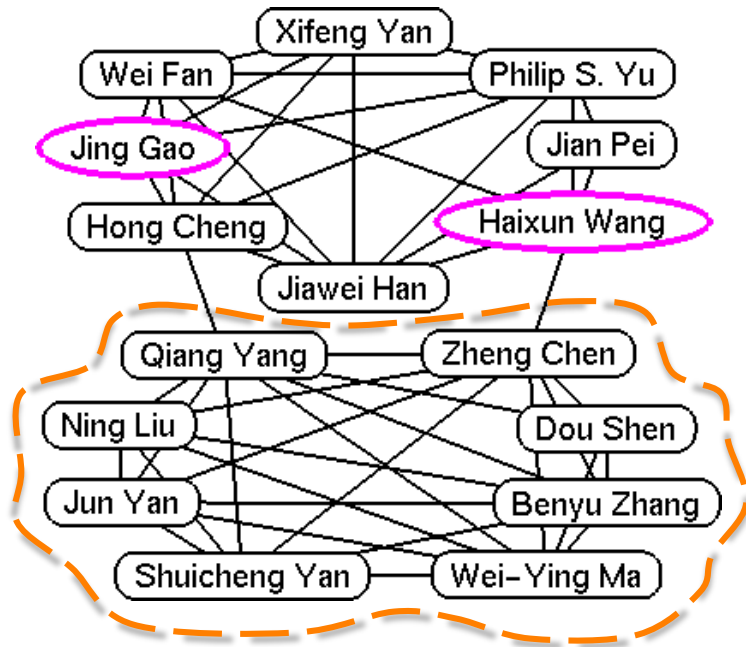
Classic density:  $|E| / |V|$

Goodness metrics	A	A U B	A U C
Classic density	2.50	<b>2.95</b>	2.83
Edge-surplus	15.3	<b>26.5</b>	22.8
Minimum degree	4	4	4
Subgraph modularity	2.0	3.6	<b>4.6</b>
Density-isolation	-2.6	<b>3.8</b>	1.5
Ext. conductance	0.25	0.14	<b>0.11</b>
Local modularity	0.63	0.70	<b>0.78</b>

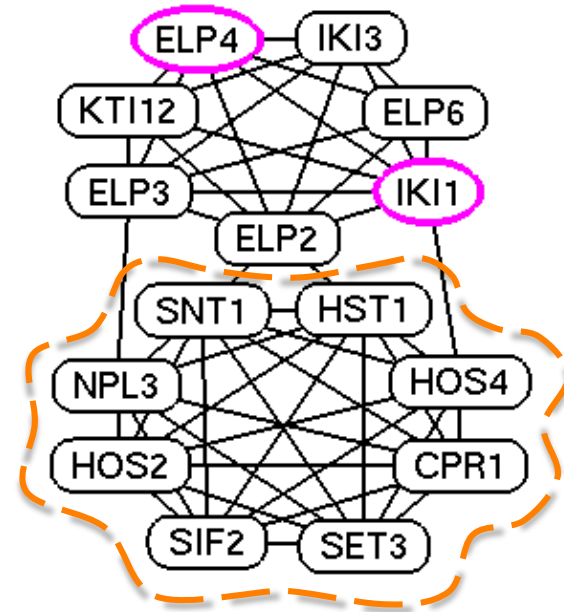
Free Riders: irrelevant to query nodes



# Free Rider Effect in Real Networks



(a) Co-author network



(b) Biological network

One existing method: classic density

# Query Biased Node Weighting

**Node Weight:**  $\pi(u) = \frac{1}{r(u)}$

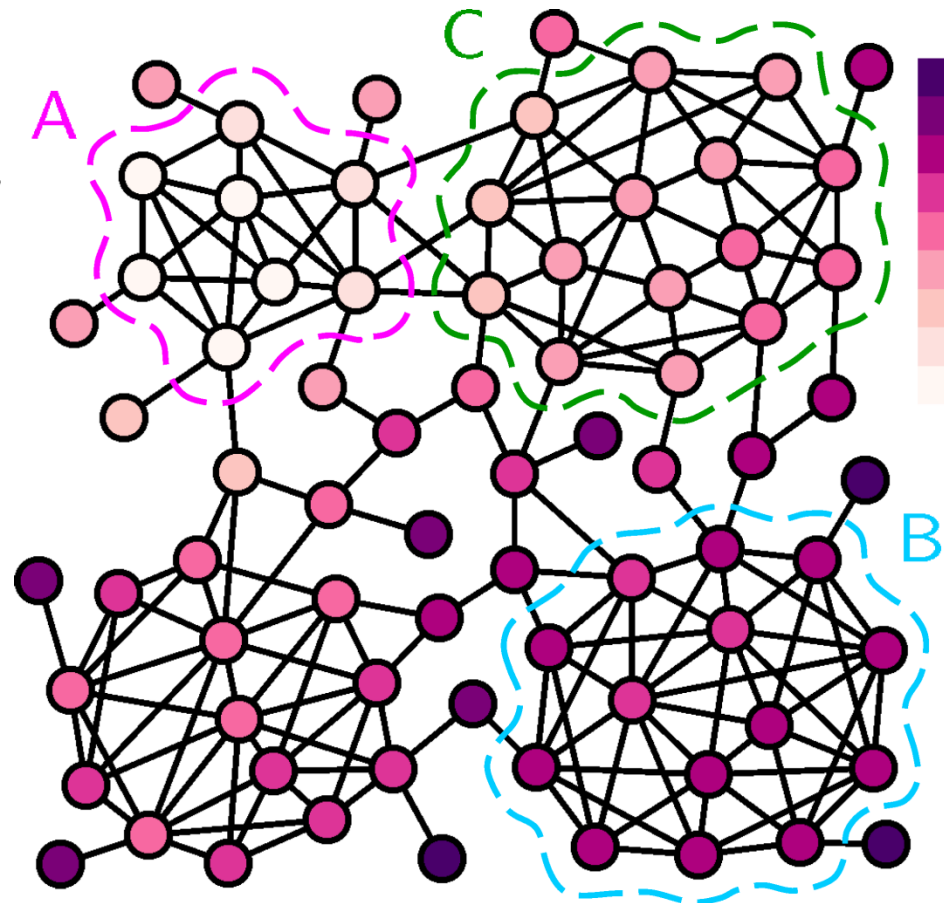
$r(u)$  : proximity value w.r.t. the query

**Query biased density:**

$$\rho(S) = \frac{e(S)}{\pi(S)}$$

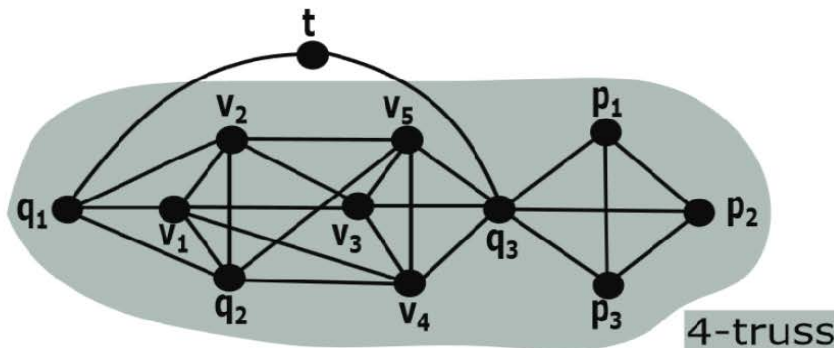
$\pi(S) = \sum_{u \in S} \pi(u)$  : sum of node weights

Subgraph **A** becomes the  
**query biased densest subgraph**

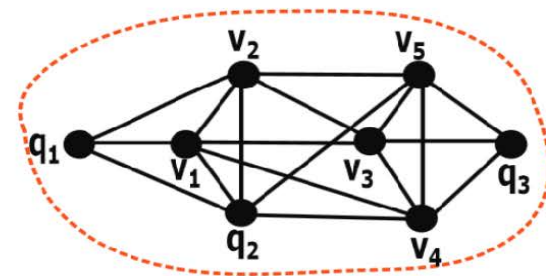


# Graph Diameter

- **Graph Diameter** of  $G$ :  $\text{diam}(G) = \max_{u,v \in G} \{\text{dist}_G(u, v)\}$
- Fig.(a), shaded, has diameter 4, the longest shortest path span from  $q_1$  to  $p_1$
- But, Fig.(b) has diameter 3.

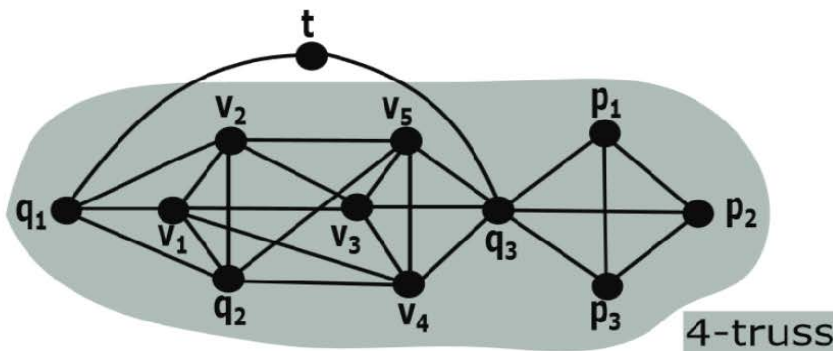
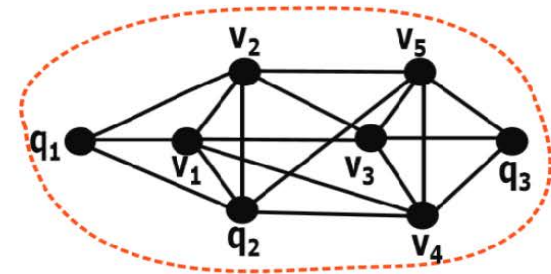


(a) Graph G

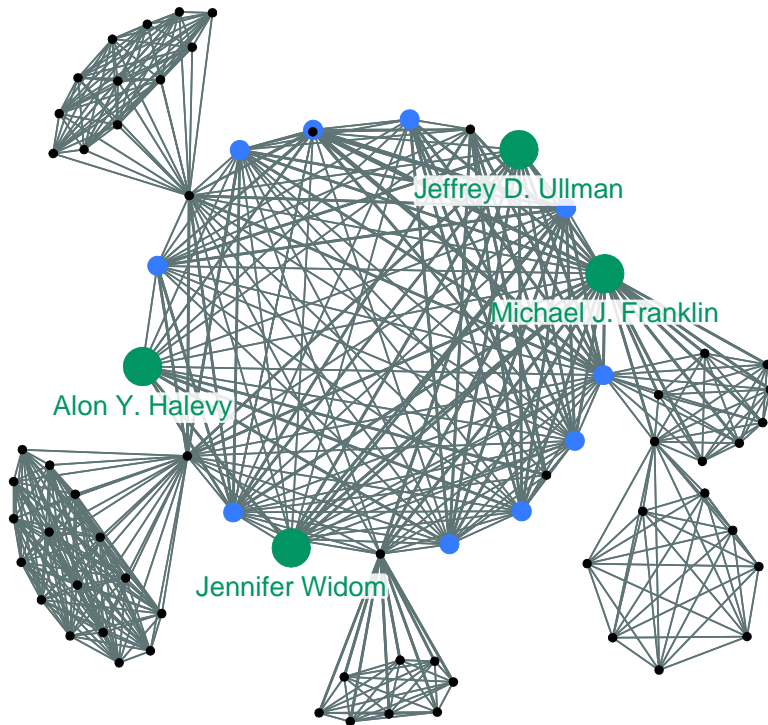
(b) Closest Truss Community for  $Q = \{q_1, q_2, q_3\}$

# Closest Truss Community Search

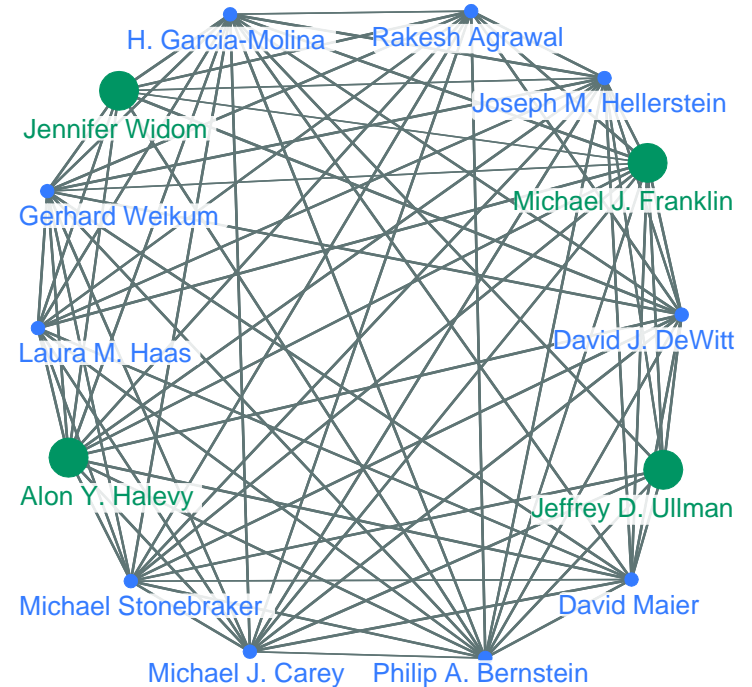
- Input:
  - a graph  $G$  & a set of query nodes  $Q$
- Output: a connected subgraph  $H$  containing  $Q$  such that
  - (1)  $H$  is a  $k$ -truss with the largest  $k$
  - (2)  $H$  has the smallest diameter among subgraphs satisfying (1).

(a) Graph  $G$ (b) Closest Truss Community for  $Q = \{q_1, q_2, q_3\}$

# Case Study: DBLP network



(a) 9-truss



(b) Closest Truss community

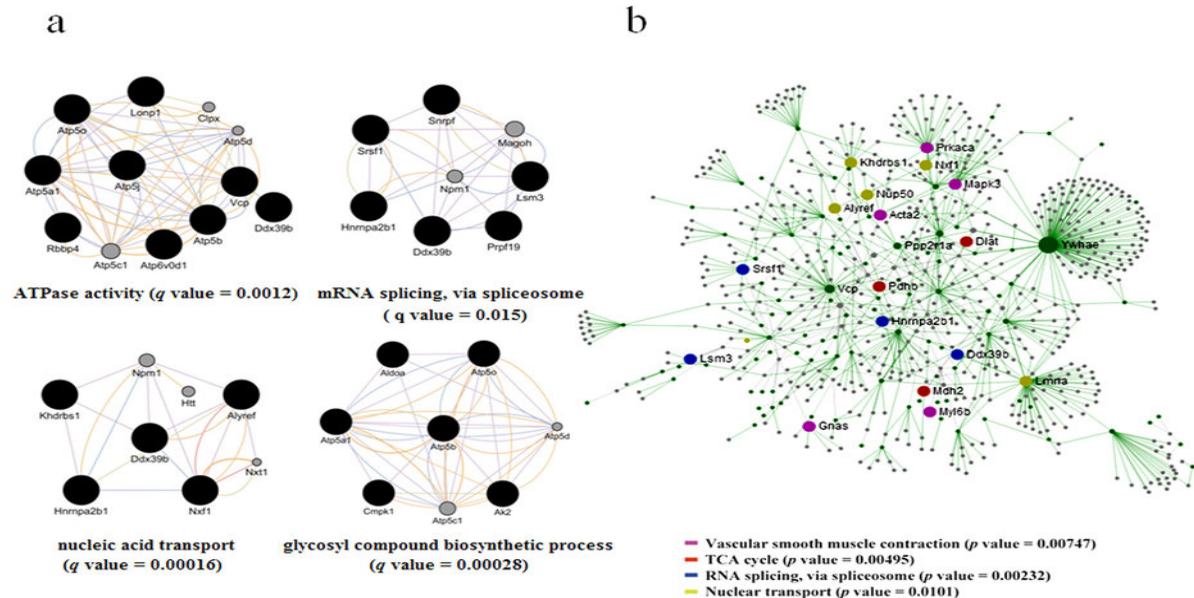
Community search on DBLP network using query  $Q = \{ \text{“Alon Y. Halevy”, “Michael J. Franklin”, “Jeffrey D. Ullman”, “Jennifer Widom”} \}$

# Desiderata of Good Query Communities

- **Query nodes:** single or multiple.
- **Cohesive structure:** quasi-clique, densest subgraph, k-core, or k-truss.
- **Quality of approximation:** guaranteed or non-guaranteed.
- **Input queries:** parameter-free or user-unfriendly.

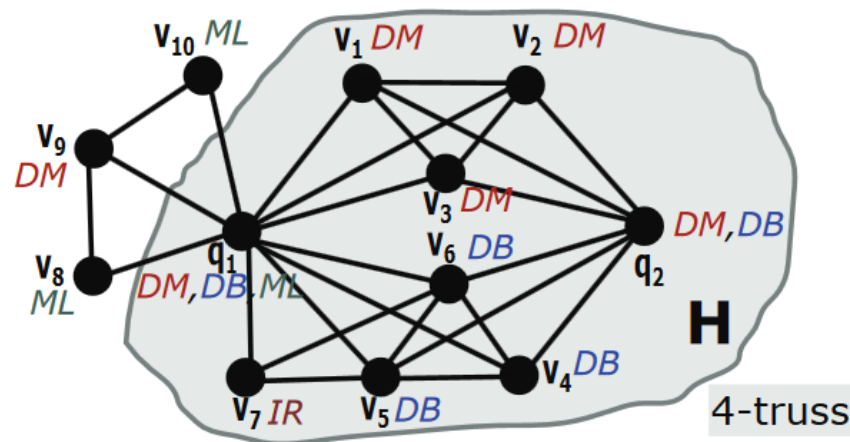
# Part 2: Attributed Community Search

- Motivation: many real social networks contain attributes or predicates on the vertices.
  - Vertices: **Person (in social networks)**, Attributes: **name, interests, and skills**.
    - Facebook: link relationship, user background
    - Twitter: following/follower-ship, tweets
  - Vertices: **Protein (in PPI networks)**, Attributes: **GO (Gene-Ontology) terms** representing **molecular functions, biological processes, and cellular components**.



# Community Search in Attributed Graph

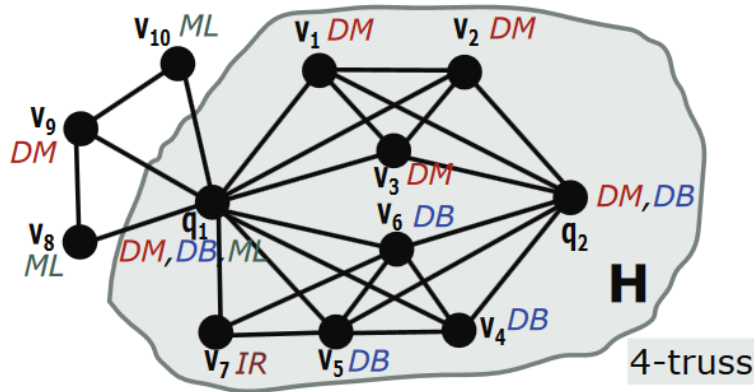
- **Structure + Semantics:** In addition to the **network structure**, users may aim to search for **attribute-related communities**, or **attributed communities**.
- Input: a **graph G** where nodes are associated with attributes  
an **input query Q** consisting of nodes  $V_q$  and attributes  $W_q$
- Output: a **connected community H** containing **Q** such that most community members are **densely inter-connected** and have **similar attributes**



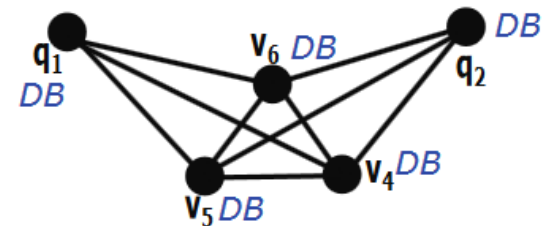
An example of  
collaboration  
attributed network



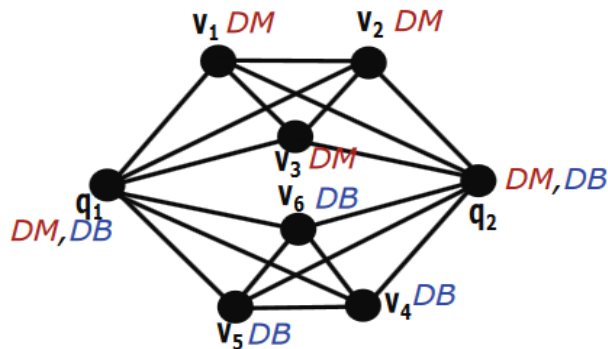
# Community Search in Attributed Graph



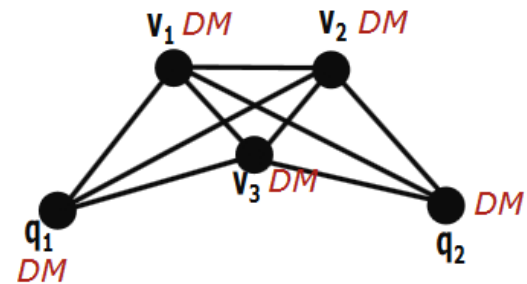
An example attributed graph  $G$



(a)  $H_1$ . 4-truss community on  $V_q = \{q_1, q_2\}, W_q = \{DB\}$



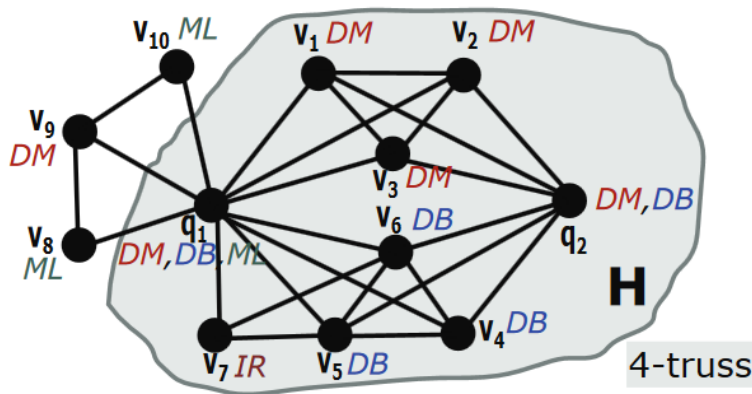
(b)  $H_2$ . 4-truss community on  $V_q = \{q_1, q_2\}, W_q = \{DB, DM\}$



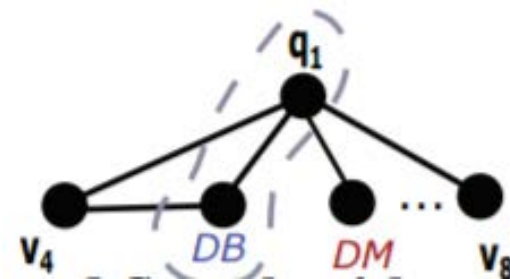
(c)  $H_3$ . 4-truss community on  $V_q = \{q_1, q_2\}, W_q = \{DM\}$

# Keyword Search

- Input: given a query consisting of nodes and attributes (keywords), e.g.,  $W=\{q_1, DB\}$
- Output: finds the substructure (trees or subgraphs) with minimum communication cost that connect the input keywords/nodes, where the communication cost is based on diameter, weight of spanning tree or steiner tree.



An example attributed graph  $G$



Keyword Search with query  $W=\{q_1, DB\}$

# A Comparison of Representative Works

- Keyword Search (KS), Team Formation (TF), Densely-connected Community Search (DCS) and Attributed Community Search (ACS)

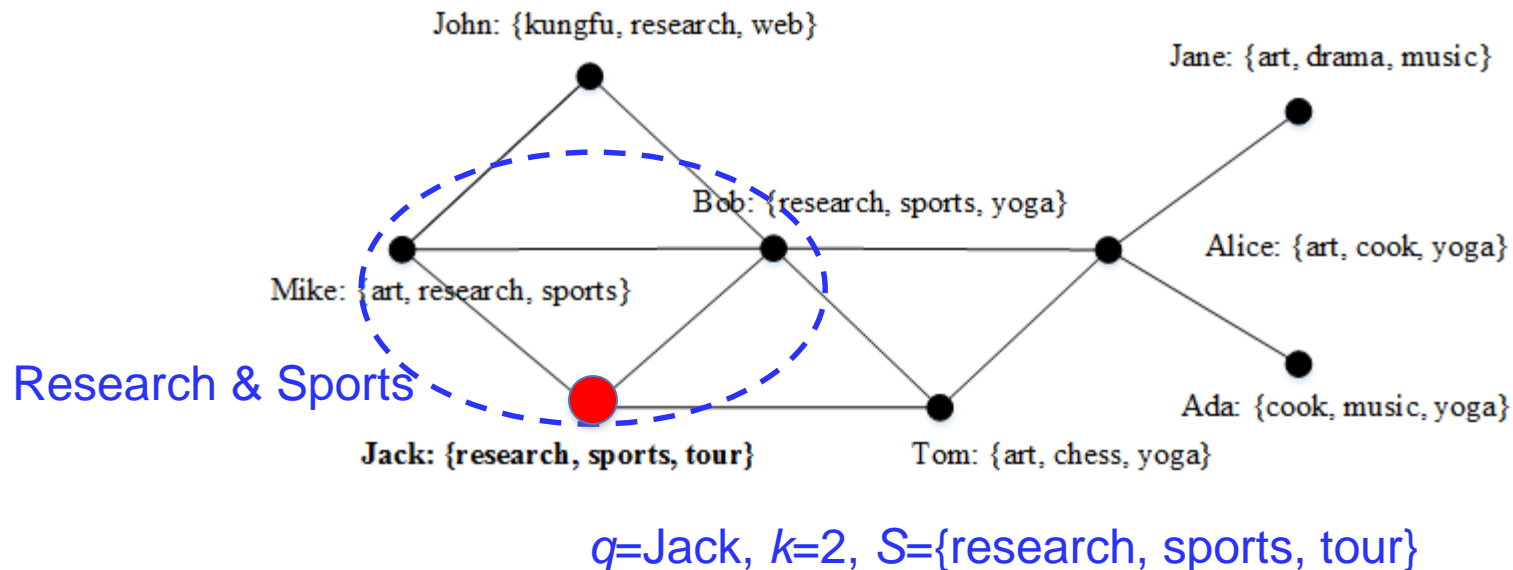
Method	Topic	Participation Condition	Attribute Function	Cohesiveness Constraint	Communication Cost
[6]	KS	✗	✓	✗	✓
[17]	KS	✗	✓	✗	✓
[30]	KS	✗	✓	✗	✓
[29]	TF	✗	✓	✗	✓
[19]	TF	✗	✓	✓	✓
[28]	TF	✗	✓	✗	✓
[39]	DCS	✓	✗	✓	✓
[14]	DCS	✓	✗	✓	✗
[15]	DCS	✓	✗	✓	✗
[5]	DCS	✓	✗	✓	✗
[26]	DCS	✓	✗	✓	✓
[31]	DCS	✗	✗	✓	✗
[46]	DCS	✓	✗	✓	✓
[18]	ACS	✓	✓	✓	✗
[25]	ACS	✓	✓	✓	✓

# The Number of Related Works

Graph type	Community Detection	Community Search
<b>Non-attributed</b>	[1000+ papers]	[10+ papers]
<b>Attributed</b>	[100+ papers]	<b>K-core-based: ACQ</b> <b>K-truss-based: ATC</b>

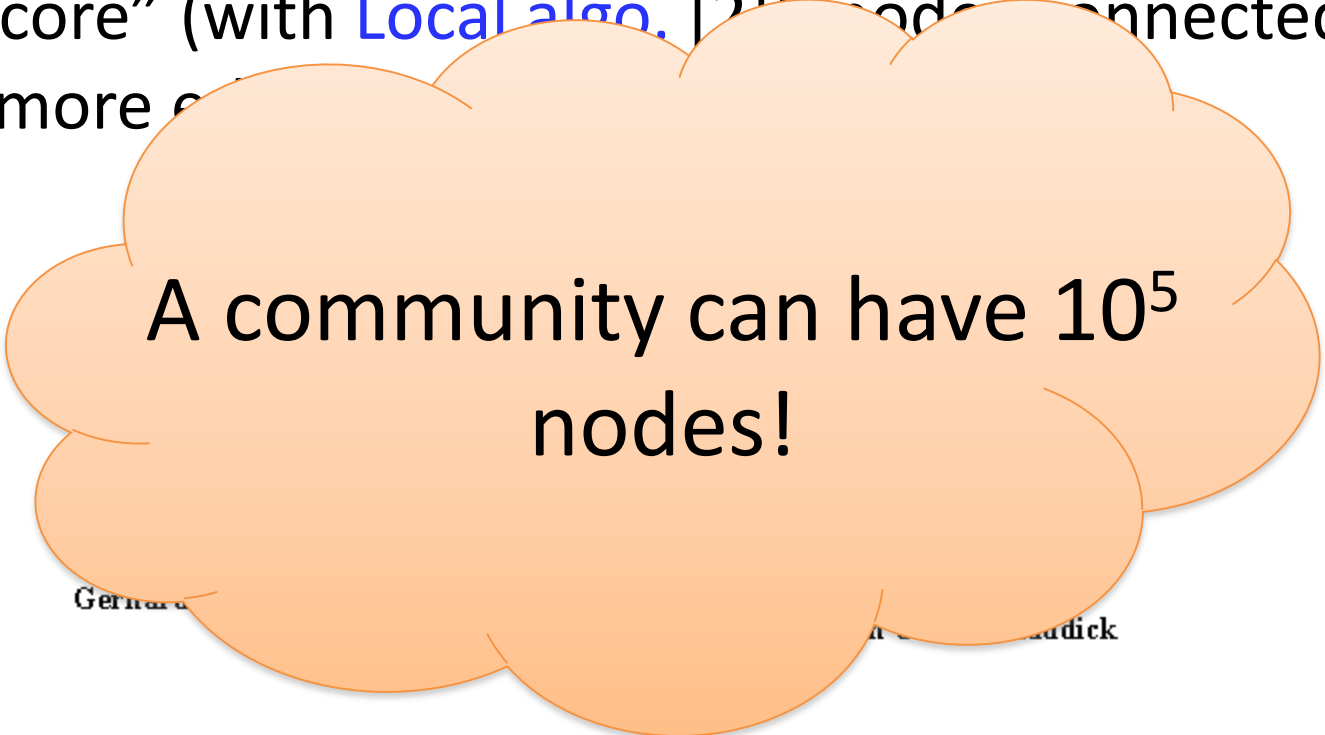
# Attributed Community Query (ACQ)

- Given a graph  $G$ , a vertex  $q$ , a set  $S$  of keywords and an integer  $k$ , find the sub-graphs s.t. each  $G_q$  satisfies:
  - Connectivity:**  $G_q$  is connected and it contains  $q$ ;
  - Structure cohesiveness:** minimum degree  $\geq k$ ;
  - Keyword cohesiveness:** the number of keywords in  $S$  shared by other vertices in  $G_q$  is maximized



# Densely-connected Community Search [1,2]

- Who is in Jim Gray's community?
  - “k-core” (with [Local algo. \[2\]](#)) nodes connected by k=4 or more edges



A community can have  $10^5$  nodes!

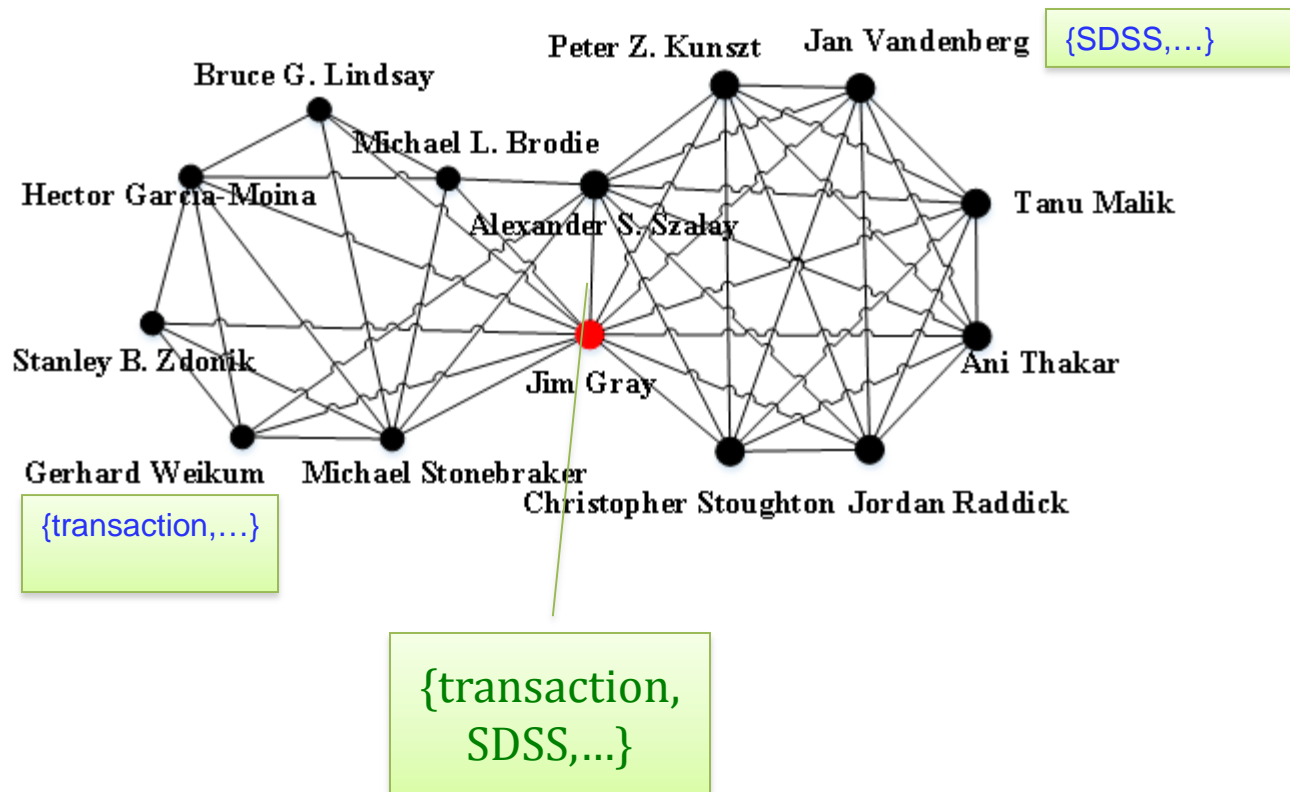
- **Why** are these people considered as Jim's community?
- What is the **theme** of this community?

[1] Sozio, Mauro, and Aristides Gionis. "The community-search problem and how to plan a successful cocktail party." *Proceedings of the 16th ACM SIGKDD international conference on Knowledge discovery and data mining*. ACM, 2010.

[2] Cui, Wanyun, et al. "Local search of communities in large graphs." *Proceedings of the 2014 ACM SIGMOD international conference on Management of data*. ACM, 2014.

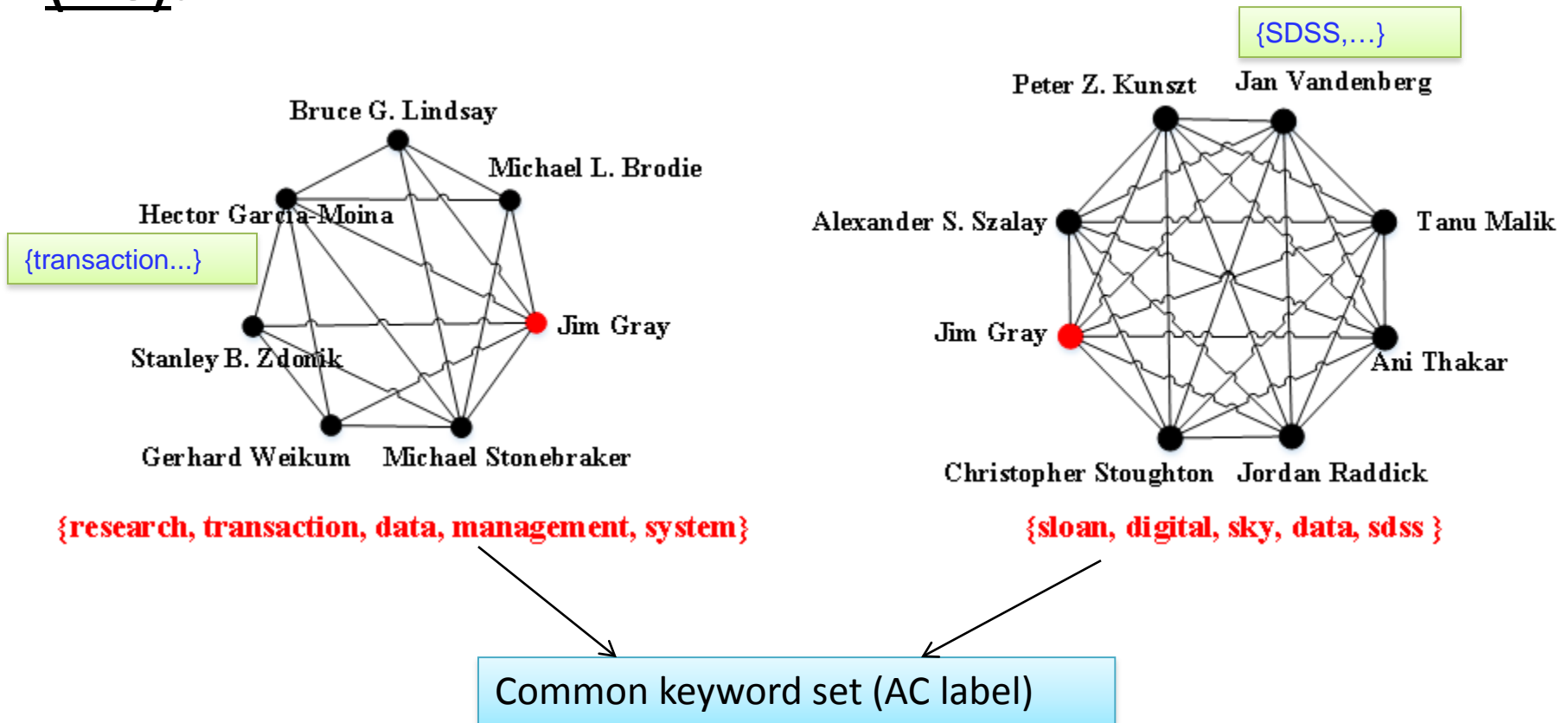
# Attributed Community (AC)

- Previous CS solutions overlook keywords
  - e.g., a researcher's interest



# Attributed Community (AC)

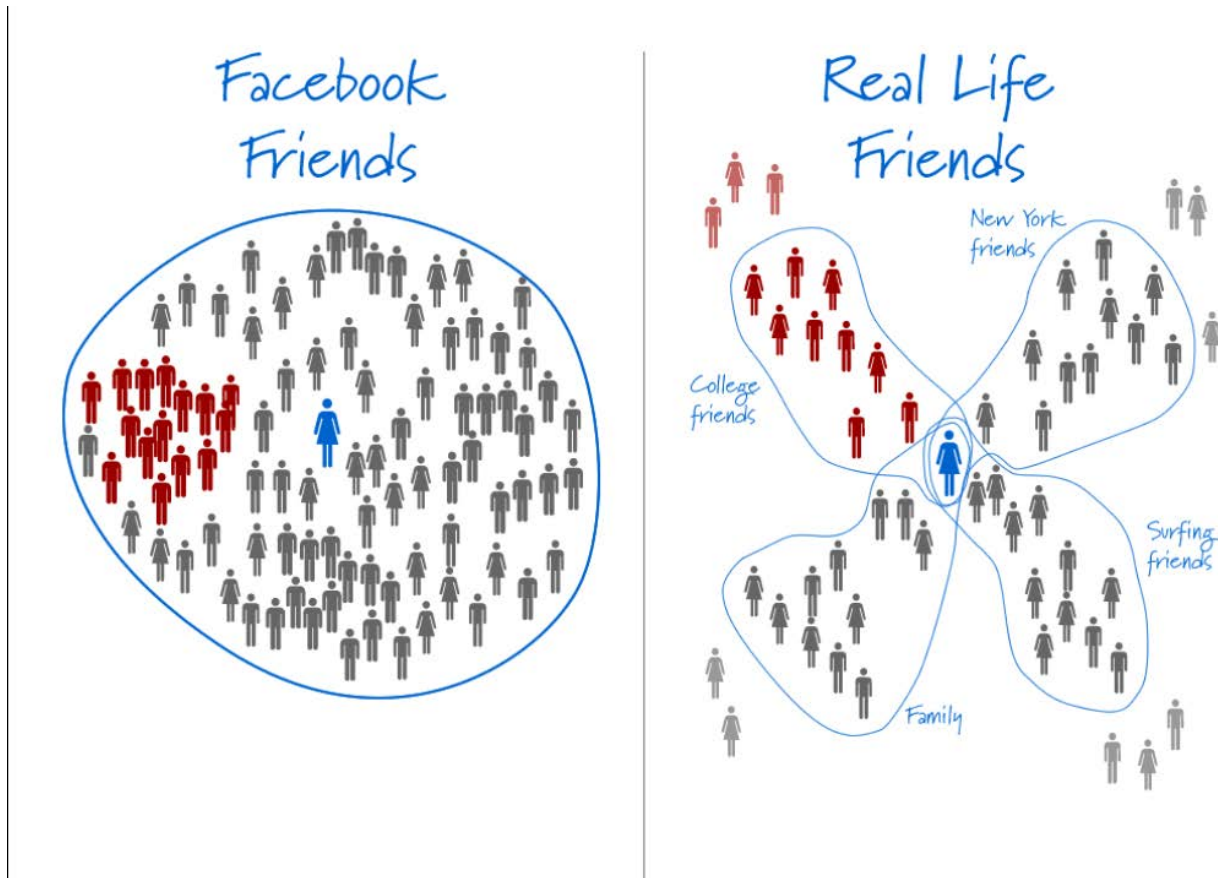
- In fact, Jim has 2 distinct attributed communities (AC).





# Part 3: Social Circle Discovery

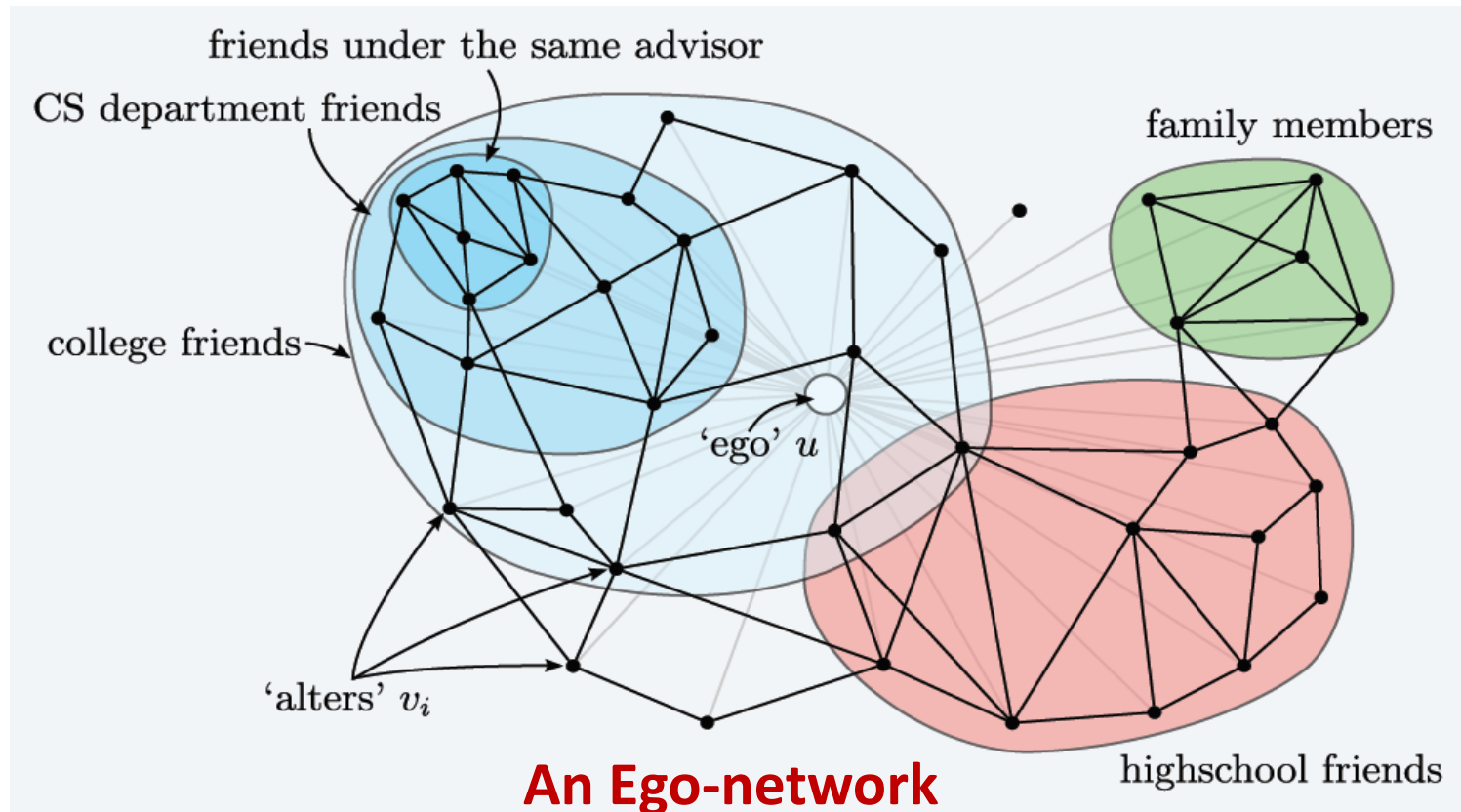
- **Social circles:** communities formed by only friends



Social Circle in Facebook

# An Ego-network

- **Ego-network**: an induced subgraph of a network only by her friends.



# Social Circle Discovery

- **Examples:** online social networks allow users to manually categorize their friends into social circles within their ego network (e.g., circles on Google+)
- **Social circle discovery:** the task is to automatically identify all social circles for a given user.
- **Applications:**
  - content filtering
  - privacy protection
  - sharing groups of users that others may wish to follow

# Learning to discover social circles

- **An unsupervised community model** predicts hard memberships to multiple, overlapping circles, using both **user profile** and **network structure**.

$$p((x, y) \in E) \propto \exp \left\{ \underbrace{\sum_{C_k \supseteq \{x, y\}} \langle \phi(x, y), \theta_k \rangle}_{\text{circles containing both nodes}} - \underbrace{\sum_{C_k \not\supseteq \{x, y\}} \alpha_k \langle \phi(x, y), \theta_k \rangle}_{\text{all other circles}} \right\}$$

Training is done by maximum likelihood, using QPBO and L-BFGS.

# Datasets: Ground-truth Social Circles

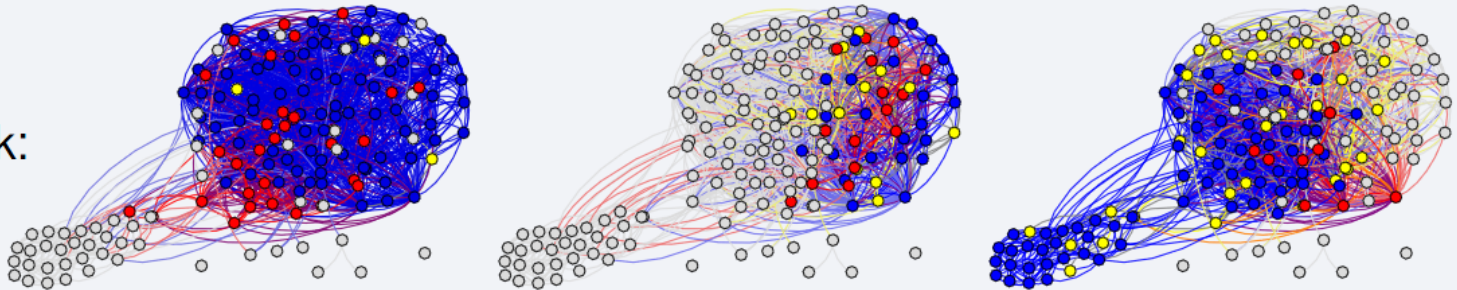
- Datasets are collected from real-world networks  
**Facebook, Google+, and Twitter**

	ego-networks	circles	nodes	edges
Facebook	10	193	4,039	88,234
Google+	133	479	107,614	13,673,453
Twitter	1,000	4,869	81,306	1,768,149

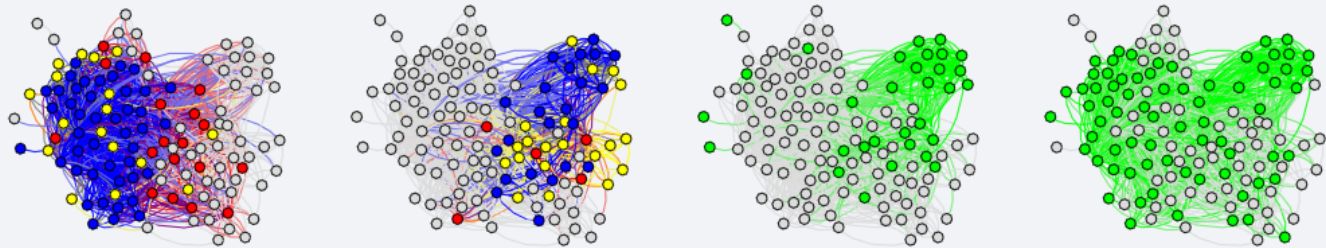
All data are **available** on [snap.stanford.edu/data/](http://snap.stanford.edu/data/)

# Detected Circles

Facebook:



Google+:



Blue = true positive; gray = true negative; red = false positive; yellow = false negative; green = detected circles for which we have no groundtruth.

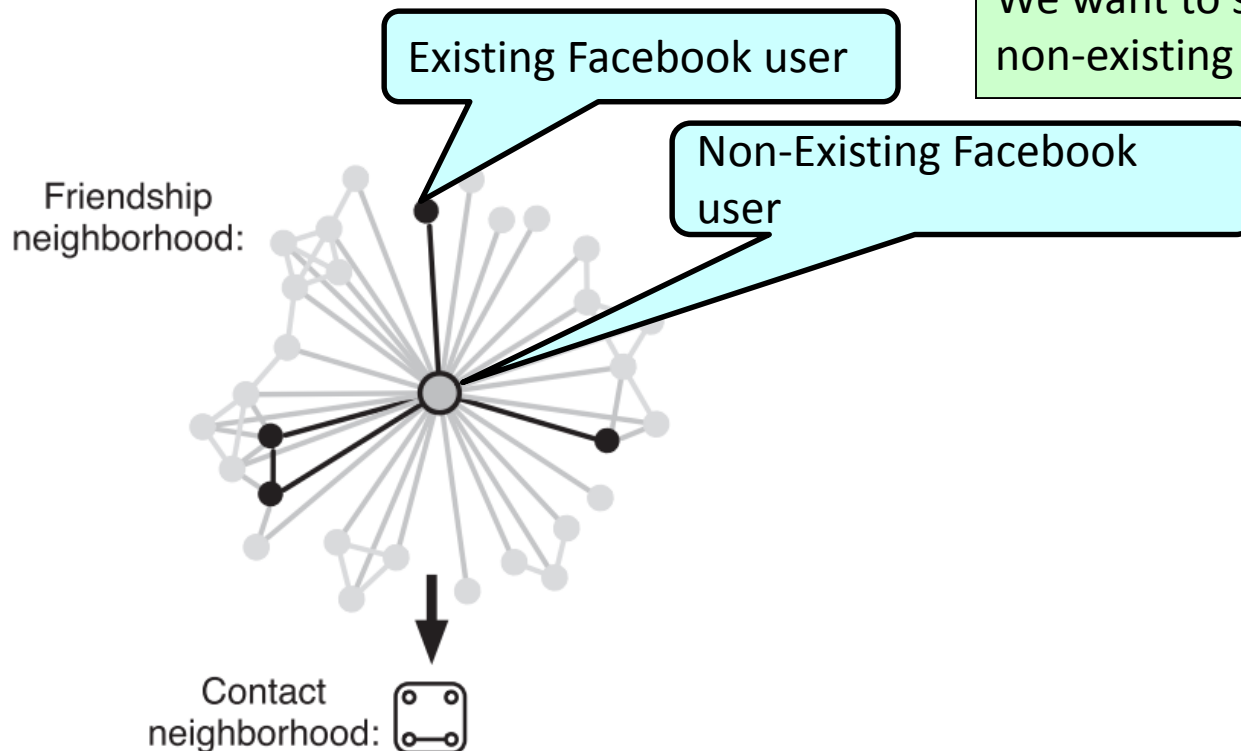
# Social Contagion

- Case Study (Facebook)  
*[Ugander et al, PNAS'12]*

Social circles can affect the process of information diffusion on social contagion

Consider an existing Facebook user invites the non-existing Facebook user to join Facebook.

We want to study the success rate that this non-existing user will join Facebook



# Social Contagion

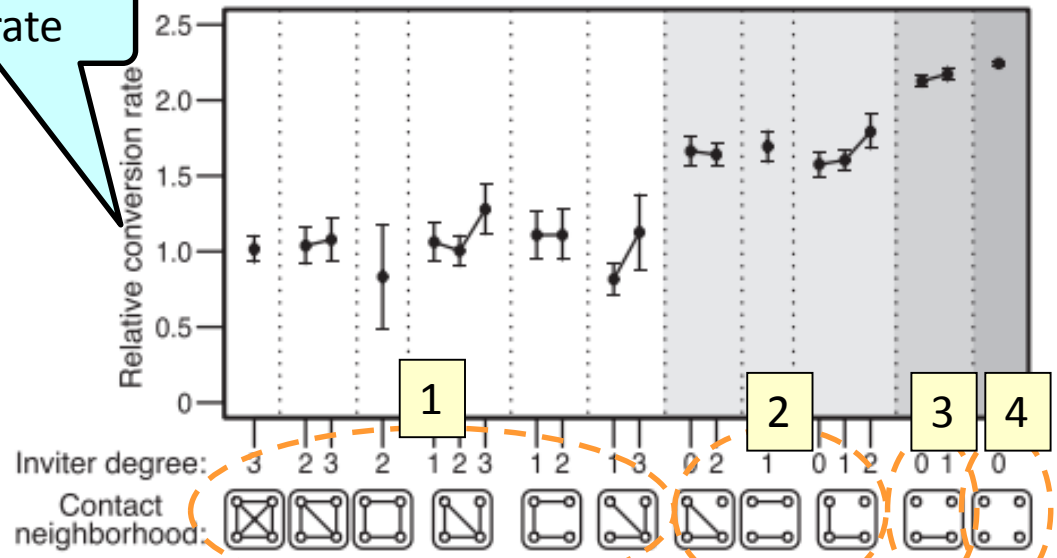
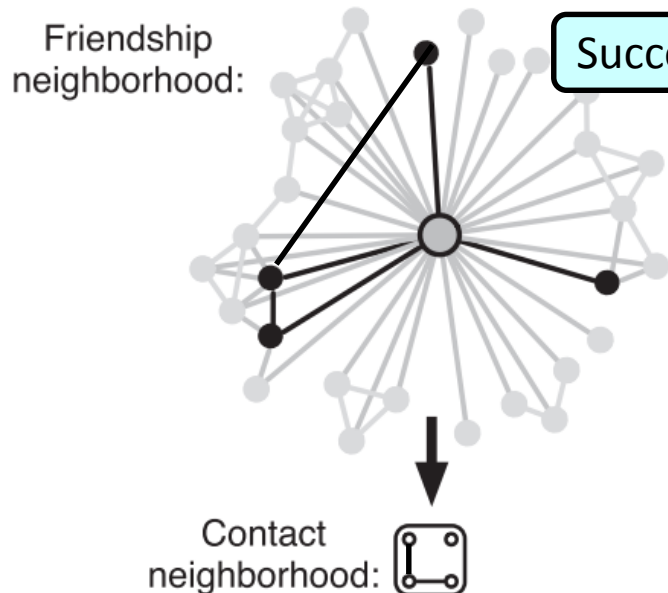
- Case Study (Facebook)

[Ugander et al, PNAS'12]

Social circles can affect the process of information diffusion on social contagion

Consider an existing Facebook user invites the non-existing Facebook user to join Facebook.

The no. of connected components is related to the success rate



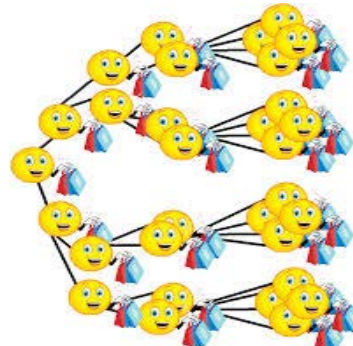


# Top-K Structural Diversity Search

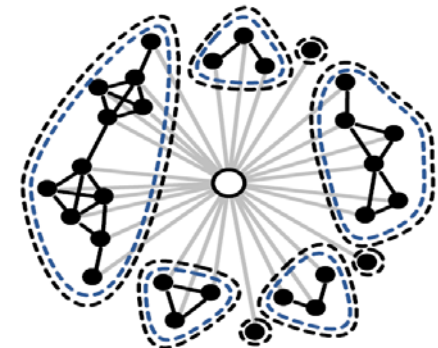
- The **structural diversity** of a node is defined to be **the number of connected components** in its ego-network.
- **Problem**
  - **Find k nodes** with the **greatest structural diversity** in a social network (Node Ranking).
- **Application**
  - Political campaign, promotion of health practices, marketing



Opinions Diffusion



Viral Marketing



Connected Components  
in the Neighborhood

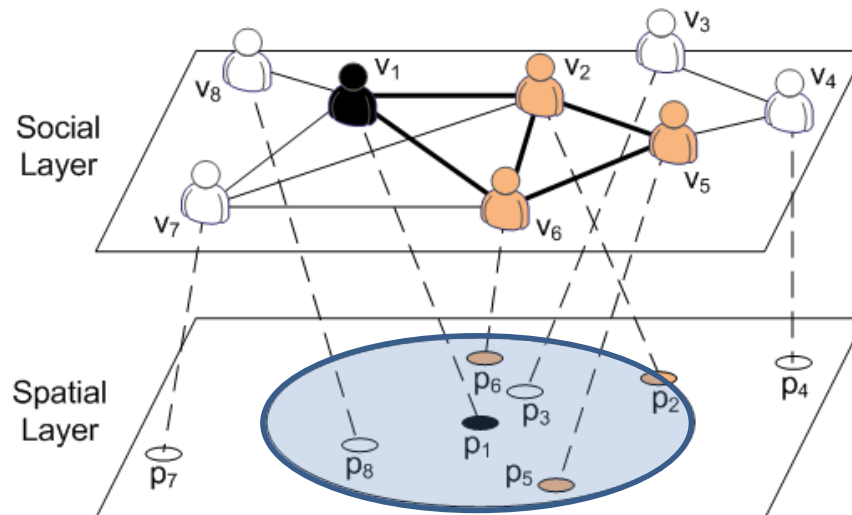
# Part 4: Querying Geo-Social Groups

- Boom in geo-social networks
  - Foursquare, Facebook, Weibo, DaZhongDianPing, Yelp, Flickr
  - Social networks coupled with user locations
- Group-based activity planning and business
  - Find a group of friends at the conference for gathering
  - Find a group of nearby friends for sports, ridesharing,groupon...



# Geo-Social Group Queries (GSGQ)

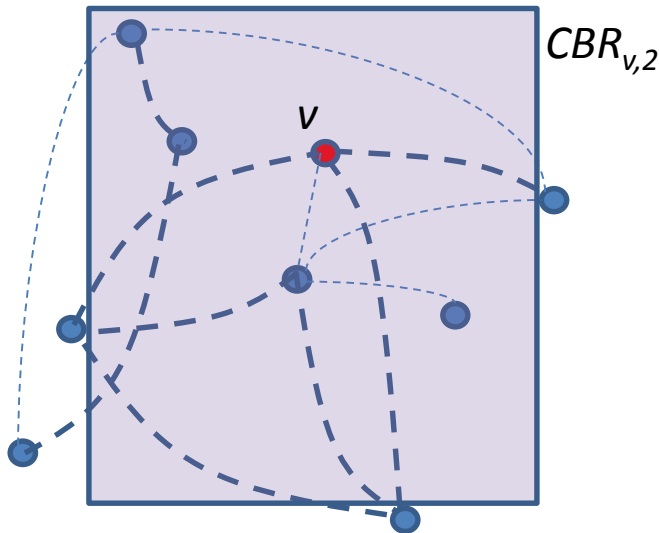
- Given an LBSN  $G=(V, E)$ , a query user  $v_q \in V$  and an integer  $c \geq 1$ , find a group of users  $V' \subseteq V$  containing  $v_q$  and satisfying:
  - Social constraint:**  $G[V'] \in G$  is a  $c$ -core
  - Spatial constraint:**
    - Range:** all users of the group are in a given spatial range
    - kNN:** the closest group with  $k$  other users (**NP-hard!**)



- Range:**  $c = 2$ ,  
 $V' = \{v_1, v_2, v_5, v_6\}$
- kNN:**  $c = 2, k = 2$   
 $V' = \{v_1, v_2, v_6\}$

# Key Concept

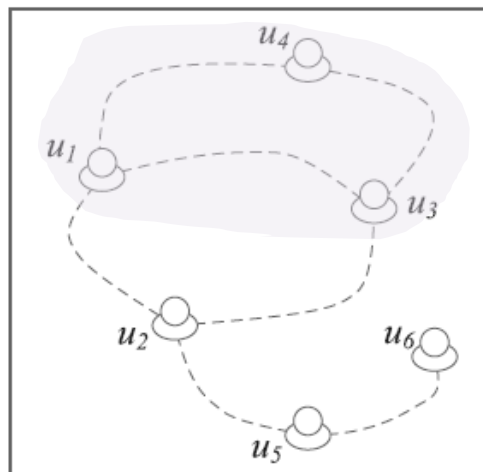
- **Core Bounding Rectangle (CBR):** Given  $G=(V, E)$ , a node  $v$ , an integer  $c \geq 1$ ,  $CBR_{v,c}$  is a rectangle that covers  $v$  and in which any user group containing  $v$  cannot form a  $c$ -core.
  - $CBR_{v,c1} \subseteq CBR_{v,c2}$ , if  $c1 < c2$
  - Construction cost:  $O(|E| \log |V|)$



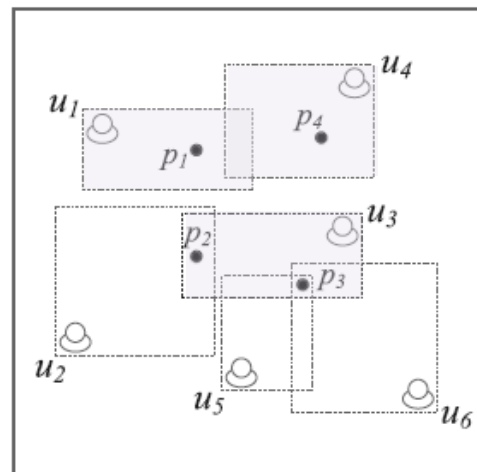
**Pruning:** exclude  $v$  from result group if query range  $\subset$  CBR

# Geo-Social K-Cover Group Queries

- Problem:** Given an LBSN  $G(V, E)$ , a set of query points  $P = \{p_1, p_2, \dots, p_m\}$ , and an integer  $k \geq 1$ , find a group of users  $V' \subseteq V$  satisfying:
  - Spatial constraint:**  $P \subset \cup_{u \in V'} u.R$
  - Social constraint:**  $G[V'] \in G$  is a  $c$ -core
  - Size requirement:**  $|V'|$  is minimum



(a) Social networks



(b) Associated regions

- $c = 2,$   
 $P = \{p_1, p_2, p_3, p_4\}$   
 $V' = \{u_1, u_3, u_4\}$

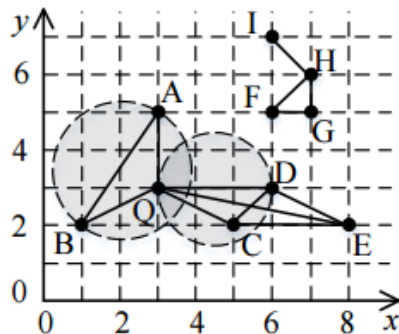
# Applications

- **Spatial task outsourcing:** identify a group of workers whose service regions collectively cover the locations of spatial tasks
- **Travel Recommendation:** find a minimum group of tourists for a self-drive tour of a set of POIs
- **Collaborative team organization:** find a collaborative team to promote products in several market areas



# Other Geo-Social Group Queries

- Spatial-Aware Community (SAC) Search
  - Y. Fang, et al., “Effective Community Search over Large Spatial Graphs” [PVLDB’17]
  - **Problem:** Given a graph  $G(V, E)$ , an integer  $c$ , and a query vertex  $q \in V$ , find a subgraph  $G_q \subseteq G$ :
    1. **Connectivity:**  $q \in G_q$  is connected
    2. **Structure cohesiveness:**  $\forall v \in G_q, \deg_{G_q}(v) \geq c$
    3. **Spatial cohesiveness:** smallest *minimum covering circle*



- $q=Q$  and  $c=2$ ,  
 $G_q = \{Q, C, D\}$

# Open Problems & Future Directions

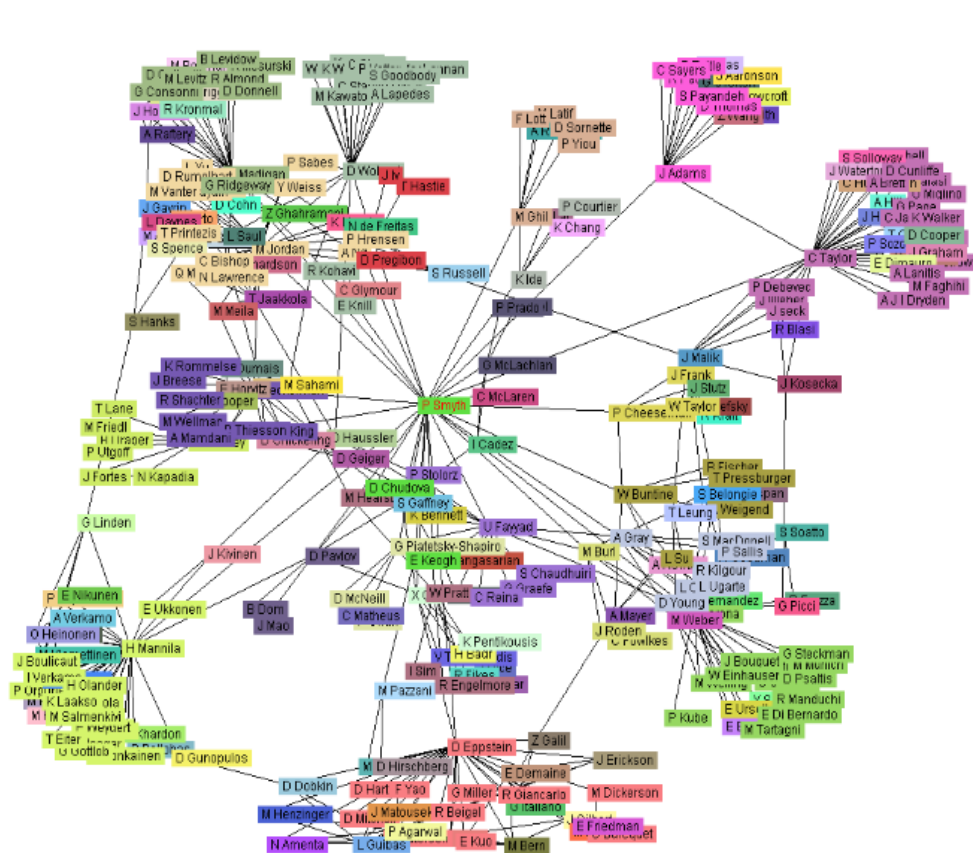
- **Heterogeneous Information Networks**
- **Scalability**
  - I/O-efficient algorithms & distributed computing
  - Stream graphs
- **Public-Private Social Networks**
- **Community Search on Uncertain Graphs**
  - Probabilistic k-core & Probabilistic k-truss



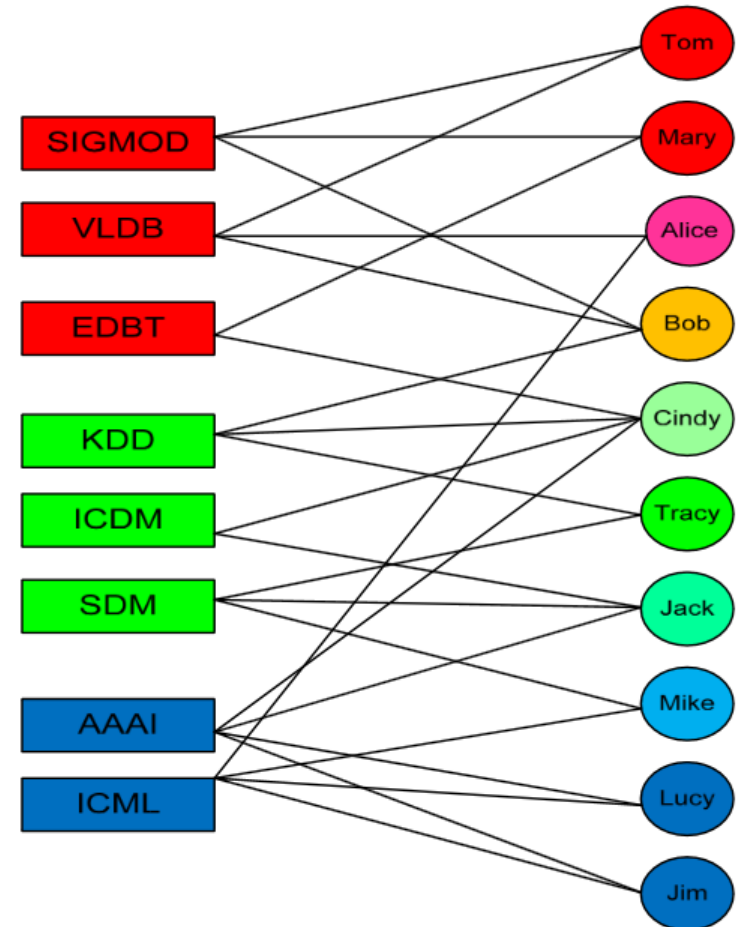
# Heterogeneous Information Networks

- Information network: A network where each node represents an entity (e.g., actor in a social network) and each link (e.g., tie) a relationship between entities.
- Homogeneous vs. heterogeneous networks
  - Homogeneous networks
    - Single object type and single link type
    - Single model social networks (e.g., friends)
  - Heterogeneous, multi-typed networks
    - Multiple object and link types
    - Healthcare network: patients, doctors, disease, hospitals, treatments

# Heterogeneous Information Networks



**Co-author Network**



**Conference-Author Network**

# Open Problems & Future Directions

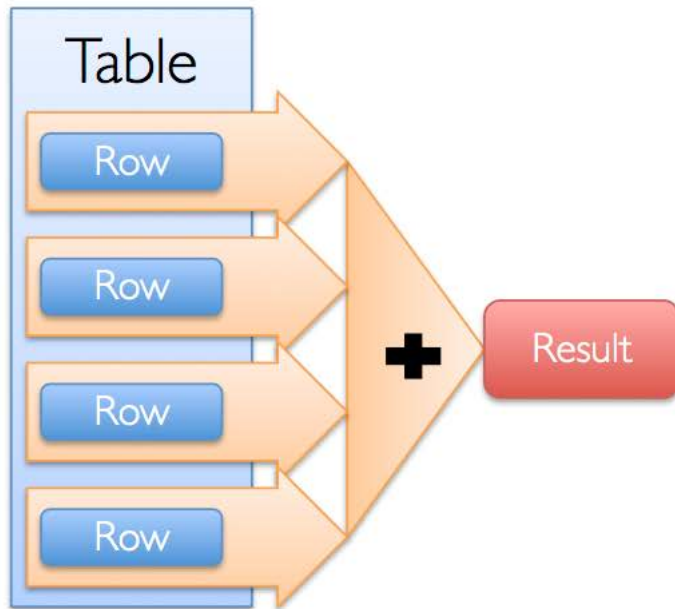
- Heterogeneous Information Networks
- **Scalability**
  - I/O-efficient algorithms & distributed computing
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# Scalability

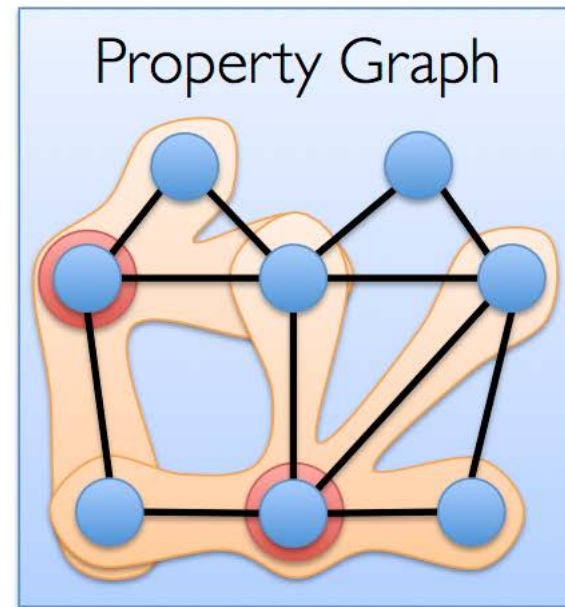
- **Scaling community search techniques** to the *massive and rapidly growing network datasets* of the Big Data era.
- **I/O efficient algorithms:** k-core decomposition and k-truss decomposition.
- **Distributed graph computing:** Pregel and Blogel.
- **Streaming graphs:** handling community indexes in highly evolving graphs.

# Scalability

Data-Parallel



Graph-Parallel



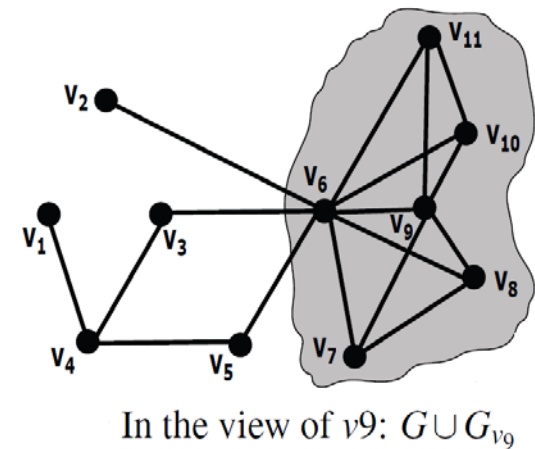
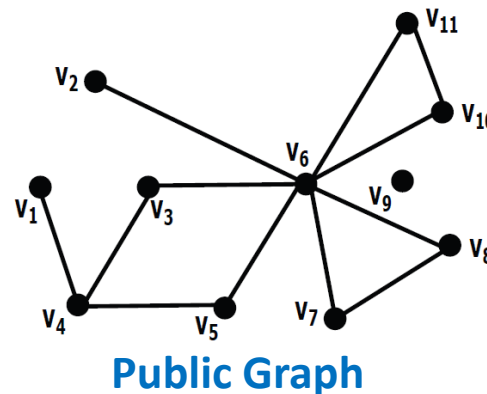
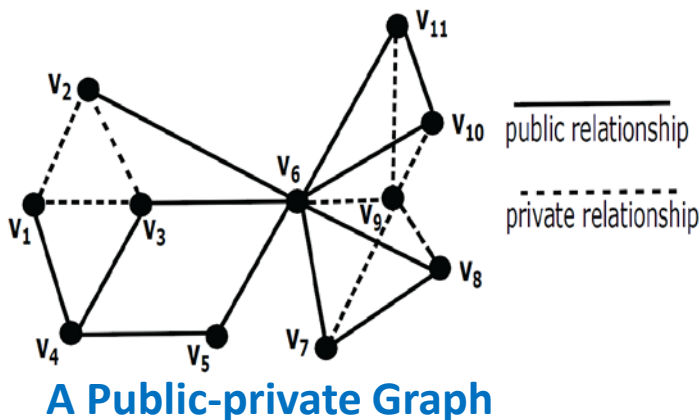
# Open Problems & Future Directions

- Heterogeneous Information Networks
- Scalability
  - I/O-efficient algorithms & distributed computing
  - Stream graphs
- **Public-Private Social Networks**
- Community Search on Uncertain Graphs
  - Probabilistic k-core & Probabilistic k-truss

# Public-Private Social Networks

- **Background:** In Facebook social network, 52.6% of 1.4 million New York City Facebook users hid their friends list.  
微博悄悄关注(Secretly follow in Weibo networks)

- Public-Private graph model contains a **public graph**, in which **each node** is also associated with a **private graph**.
  - The **public graph** is **visible to everyone**, but each **private graph** is **visible only to the corresponding user**.



# Open Problems & Future Directions

- Heterogeneous Information Networks
- Scalability
  - I/O-efficient algorithms & distributed computing
  - Stream graphs
- Public-Private Social Networks
- **Community Search on Uncertain Graphs**
  - Probabilistic k-core & Probabilistic k-truss

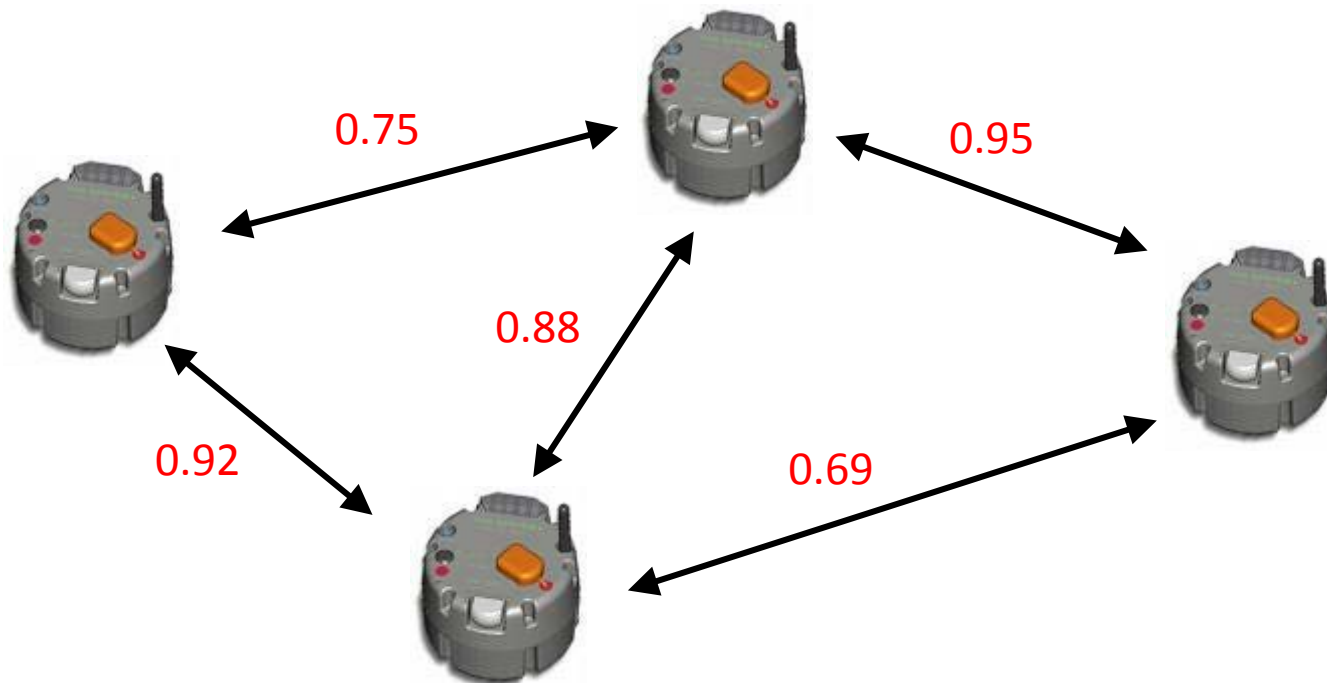


Not all real-world networks  
are  
deterministic graphs.

Probabilistic/Uncertain Graphs: each edge has  
an **existence probability**.

# Probabilistic Graphs: Examples

- Topologies of wireless sensor networks (WSNs)
  - Vertices: sensor nodes
  - Edges: wireless links between sensor nodes
  - **Uncertainties: probabilities of wireless links functioning**



# Discovery of communities in uncertain graphs

- Benefits:
  - Find most influential communities in social networks.
  - Functional module identification for helping critical clinical diagnosis of diseases such as cancer in biology.
- K-core and k-truss have been studied in probabilistic graphs.
- **An exciting question** is how to generalize various community models and search techniques to probabilistic graphs.

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Thank you!

Questions?

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